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A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

THESIS Donald Arthur Parish Captain, USAF

AFIT/GOR/ENS/94M-10



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A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

THESIS

Presented to the Faculty of the Graduate School of Engineering
of the Air Force Institute of Technology
Air University
In Partial Fulfillment of the
Requirements for the Degree of
Master of Science in Operations Research

Donald Arthur Parish, B.S.

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Donald Arthur Parish

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Abstract

Satellite range scheduling involves scheduling satellite supports in which a satellite and a specific remote tracking station communicate with each other within a specified time window. As the number of satellite supports continue to increase, more pressure is placed on the current manual system to generate schedules efficiently. Previous research efforts focused on heuristic and mixed-integer programming approaches which may not produce the best results. The objective of this research was to determine if a genetic algorithm approach to automating the generation of 24 hour schedules was competitive with other methods. The goal was to schedule as many supports as possible without conflict.

The genetic algorithm approach attempted to find the best priority ordering of support requests, and then used a schedule builder program to build schedules based on simple rules. A schedule was produced for seven days of representative satellite range data with slightly better results compared to earlier results using a mixed-integer programming formulation. Based on the reported results, the genetic algorithm approach presented in this research appears to be a competitive approach for generating 24-hour satellite range schedules.

A GENETIC ALGORITHM APPROACH TO AUTOMATING SATELLITE RANGE SCHEDULING

I. Introduction

Scheduling is the allocation of resources over time to perform a collection of tasks (2:2). Many important scheduling problems exist that are of interest to the Air Force. Examples include allocation of test range resources, airlift scheduling problems, and satellite communication support scheduling. Many standard scheduling problems may be formulated as mixed-integer programming problems where some decision variables take on integer values while other variables, such as time, have continuous values. While a mathematical programming formulation generally guarantees an optimal, or best possible, solution, the computation times required to find exact optimal solutions may be prohibitive for practical-sized problems. This is because many scheduling problems belong to the class of NP-complete problems (6:4).

In such NP-complete problems, the solution time may increase exponentially with the number of variables. For example, a problem with 10 times the number of variables as an original problem might take an order of 2¹⁰ times longer to solve. Because of these possible long solution times, heuristic techniques are often used to find good solutions in a reasonable amount of time. Heuristics are procedures which do not guarantee optimal solutions. They usually, but not always, provide feasible solutions, often require human ingenuity in their development, and are often quite problem-specific.

Heuristic methods for solving scheduling problems that are less dependent on the specific problem could be useful in solving problems of Air Force interest. This would be especially true when the constraints of the problem are difficult to put in mathematical form. One such method for finding good solutions to scheduling and other optimization problems is called a *genetic algorithm*.

Genetic algorithms (GAs) are artificial intelligence search methods based on the idea of natural selection and evolution. Initially developed by John Holland at the University

of Michigan (10), applications include problems in optimization and machine learning. Although initially applied to function optimization problems, genetic algorithms have also been applied to scheduling and combinatorial optimization problems. The main strength of genetic algorithms is their ability to quickly explore a large space of possible solutions for good, if not optimal, solutions. They differ from many traditional algorithms as their search is based only on the overall evaluation of a set of parameters. As such, they do not need to rely on derivative information to proceed. Genetic algorithms are also able to find solutions where multiple optimal solutions exist (7:2-5).

Problem Description

One particular scheduling problem of interest to the Air Force is the Satellite Range Scheduling (SRS) problem. The Air Force Satellite Control Network (AFSCN) must schedule over 300 command and control communications per day between nine remote tracking stations (RTSs) and approximately 100 satellites (8:1-3). Each communication between a satellite and a RTS is called a support. The satellite supports are necessary to maintain command and control, tracking, and system tests of the satellites (11:10). The scheduling of these supports must take place in a time window because each RTS is geographically separated from the others, and can only "see" each satellite for a limited time during its orbit. The time window is shorter for low-altitude satellites than for medium- or high-altitude satellites, as the low-altitude satellites pass out of the line-of-sight of the RTSs much more quickly than higher-altitude satellites. The longer time window of the higher-altitude satellites makes scheduling them less difficult than scheduling the low-altitude satellites.

Each Mission Control Complex (MCC) is responsible for the health, status, and orbital control of a subset of the total satellites. The composition of the subset depends on the mission type of the satellite. A MCC determines the length and required time windows for each support request for their satellites, as well as which RTSs can serve each request. These time windows may be more restricted than the physical visibility limits due to mission scheduling requirements. For example, a satellite may need a communication every hour after an operation is performed.

All requests from the individual MCCs are passed on to the range schedulers who are responsible for making an overall schedule. This schedule must ensure that each RTS antenna supports only one satellite at a time. The schedulers must also allow for a required turn-around time between supports to allow the RTS antennas to be reoriented. Down-time for RTS maintenance is also necessary. The resulting schedule is called the *initial* 24-hour schedule. The development of this schedule using genetic algorithms is the focus of this research.

The goal of the initial 24-hour scheduling process is to schedule as many communication supports as possible in a 24-hour period while satisfying the constraints of time windows, turnaround time, and scheduled maintenance down-time for the RTSs. If a conflict cannot be resolved by the range schedulers, they must de-conflict it with the MCCs and RTSs involved. If a support cannot be scheduled initially, it may be possible to facilitate its scheduling by changing the support requirements, by decreasing the turn-around time, or by altering other constraints. In the worst case, a support may not be scheduled. A good scheduling process minimizes the need for this deconfliction process. Although the current manual scheduling process schedules approximately 95-98% of the requested supports (8:5-2), the process is time-consuming, and could be streamlined by automating the development of the 24-hour schedule.

Previous Solution Efforts

Two recent thesis efforts have investigated automation of the 24-hour scheduling process. Gooley used both a mixed-integer programming (MIP) approach and heuristic scheduling methods to schedule satellite supports (8). As a follow-on research effort, Schalck improved Gooley's solution by reducing the number of integer variables necessary to define the problem (14). In both cases, the problem was too large to be solved in its entirety and had to be decomposed into smaller problems to find a solution in reasonable time.

Computational Complexity

Computational problems are often classified based on their complexity. According to Garey and Johnson (6:4), problems that are classified as NP-complete can take an inordinate amount of time to solve if the size of the problem is large enough. The general resource-constrained scheduling problem is in this class. According to Gooley, the SRS problem is a type of resource-constrained scheduling problem (8:2-9), and a fast (polynomial) solution cannot be guaranteed. Genetic algorithms, which have had some success in finding good solutions to other NP-complete scheduling problems, may be useful when applied to the SRS problem.

The mixed-integer programming approach can find optimal solutions for the satellite range scheduling problem when the number of support requests is small. However, when a larger problem is decomposed into smaller subproblems, a global optimal solution may no longer be guaranteed. A genetic algorithm approach may be expected to produce schedules that are at least as good while meeting the requirement of a short solution time since it would attempt to generate a solution to the entire problem. Also, a genetic algorithm approach can be more flexible in handling special case scheduling requests. These issues are addressed in this research.

Research Objective

The overall research objective is to determine whether a genetic algorithm-based solution methodology can effectively be applied to the satellite range scheduling problem. Successful solutions must:

- Generate feasible 24-hour schedules which schedule the greatest number of support requests possible.
- 2. Find solutions quickly. A short solution time is useful in rescheduling as requirements change.

A secondary objective is to explore the scheduling of special-case requirements. These include supports which require simultaneous support on multiple RTSs and those that must be scheduled at fixed intervals. Rescheduling of satellite supports is also a consideration.

Overview

Chapter II is a summary of current literature relevant to genetic algorithm applications to scheduling problems and includes a review of previous work on the satellite range scheduling problem. Chapter III explains the methodology used in constructing a genetic algorithm solution to the satellite range scheduling process and details the implementation of the solution. Chapter IV presents the results and compares them to past efforts. Finally, conclusions and recommendations for future work are presented in Chapter V.

II. Literature Review

This chapter summarizes current literature on genetic algorithms (GAs) as related to scheduling problems. It begins with the development of genetic algorithms. Next, the extension of genetic algorithms to combinatorial optimization problems is reviewed. Discussion then turns to applications of genetic algorithms in scheduling problems, especially to problems with similarities to the satellite range scheduling problem. Finally, past research efforts in satellite range scheduling are discussed.

Genetic Algorithms

Introduction. Genetic algorithms were developed by John Holland as part of his artificial intelligence research on how artificial adaptive systems can evolve, or change, in response to their environment in order to solve problems (10). His ideas were based on the biological theory of evolution where variations in chromosomes, or genetic codes, result in different traits of the individual. These traits of an individual in turn result in a level of performance, or fitness, of an individual. Usually, the more fit individuals in a population survive from one generation to another and reproduce. Sexual reproduction by two individuals produces individuals with new chromosomes formed by a combination of the chromosomes of its parents. If a child inherits good parts of chromosomes from each of its parents, it has a higher level of performance than the parents. Through "survival of the fittest" and reproduction of high-fitness individuals in each generation, the population as a whole tends to evolve towards higher levels of fitness. The artificial version of this process is called a genetic algorithm. It does not seek to exactly simulate biological evolution, but the concepts of biological evolution are used to find good solutions to difficult problems.

Simple Genetic Algorithm. Most genetic algorithm work is based on the simple genetic algorithm as developed by Holland (10) and described by Michalewicz (13). Figure 1 illustrates the process and a brief review follows.

Coding. The first, and generally most difficult, step of any genetic algorithm is to choose a proper coding to map the problem solution space into a genetic string,

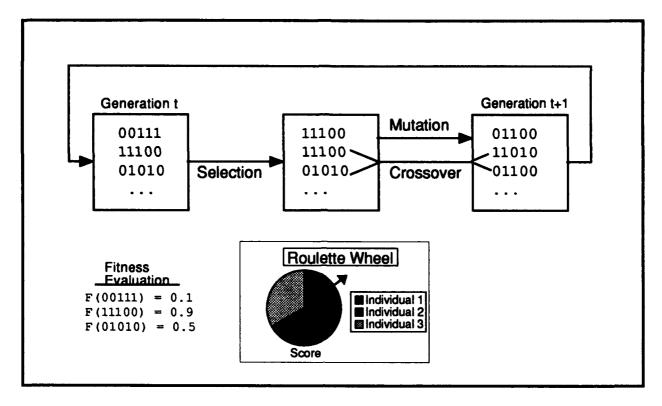


Figure 1. Genetic Algorithm Procedure

or chromosome, and to randomly create an initial population of individuals with varying strings. In the simple genetic algorithm, this coding is a binary string of zeros or ones. For function optimization, groups of binary digits are mapped so as to translate to a real number parameter representing the function value (13:19-20).

Evaluation and Selection. The strings in the population can be evaluated for their fitness relative to other strings in the population by entering their parameters into an evaluation or fitness function. The best strings reproduce by mating with each other to produce offspring for the next generation of the population. In the simple genetic algorithm, selection is governed by a "roulette wheel" selection operator. Each string has a probability of reproducing in proportion to the ratio of its fitness and the total fitness of the population. As in Figure 1, these ratios can be shown as pieces of a circular pie. New strings are selected for reproduction by randomly "spinning" the wheel. The strings whose proportion of the pie are greatest should, on average, be selected more often than the less-fit strings.

Other selection rules can be used (13:62). These include selection by ranking, where instead of a continuous scale based on fitness the strings are ordered, or ranked, by their fitness. The higher-ranked strings are given a greater chance of reproducing than average or low-ranked strings. In this way, the more fit individuals tend to reproduce.

Crossover Operator. During mating, two strings swap part of their "genetic" material. The children of these matings have parts of each of their parent's string. The swapping of genetic material, called crossover, allows for new child strings to be created, combining good aspects of their parents. The resulting strings with above-average fitness tend to survive and prosper, while those with below-average fitness tend to die out. An example of crossover is shown in Figure 2.

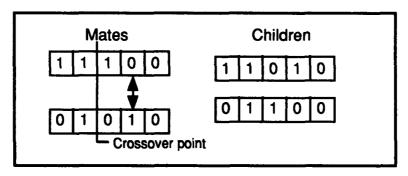


Figure 2. Example of Crossover

Mutation. Mutation randomly changes part of the string of a child to help maintain a diverse population. It is not as important as crossover in some applications since it merely acts as a type of random search (7:14).

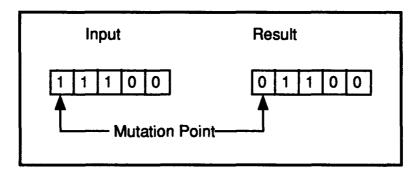


Figure 3. Example of Mutation

The general genetic algorithm thus consists of three steps which are repeated for each generation: evaluation/selection, crossover and mutation. These steps continue until termination conditions, such as some predetermined number of generations, are met. Often the genetic algorithm ends when the population converges: all strings evaluate to the same fitness (13:56). This string or strings may be the best answer, but a genetic algorithm may often converge to a suboptimum, in which case the population has prematurely converged. The major variables controlled by the experimenter to combat premature convergence are: population size, type and probability of crossover, type and probability of mutation, and selection operators (7:106-124).

Population Size. A genetic algorithm population contains a fixed number of strings. This number is called the population size. The population size affects the convergence rate of a genetic algorithm by controlling the variety of genes in the population. A smaller population may converge quickly, but usually to a sub-optimum. A larger population converges more slowly, and usually, but not always, finds a better final answer.

Although the basic procedure is simple, the framework of a simple genetic algorithm has been successful in solving a wide range of problems in function optimization and has been extended to other types of problems such as combinatorial optimization (13:165,193).

Why GAs Work. Although there is no complete formal theory to explain the operation of genetic algorithms, several hypotheses which partially explain their power have been advanced (13:51). The point they make is that in each generation, the genetic algorithm combines good partial solutions in the genes of parent chromosomes to find even better solutions in child chromosomes. The following discussion briefly describes these hypotheses. A more formal treatment of genetic algorithm theory is available in Goldberg's text (7).

Schemata and Schema Theorem. A schema is a pattern of values in a gene with the alphabet 1,0,*, where "*" is the "don't care" symbol matching any position. For a chromosome to include a particular schema it must match the schema values. For example, the chromosome (01101) contains 32 schemata. These include: (0**01), (*1****), (01101),

(*****), (0*1*1). The defining length of a schema is the distance between the outermost non-* symbols. The order of a schema is the number of non-* symbols contained in the schema.

Holland's schema theorem was the first rigorous explanation of how a simple genetic algorithm works. The schema can be thought of as representing the partial solutions in chromosomes and Holland concluded that genetic algorithms manipulate schemata when they execute (13:91). In the simple genetic algorithm, individuals reproduce and increase in proportion to their fitness. The schema theorem asserts that because of this, schema associated with individuals with above average fitness tend to increase exponentially, while those schema associated with below average performance tend to occur less often in succeeding generations (7:32-33).

Building Block Hypothesis. Much of the power of genetic algorithms comes from finding good building blocks (7:41). Building blocks are highly fit, low-order schema of short defining length. Because of their short length, these blocks tend to survive, even under the disruption caused by crossover. "In a way, by working with these particular schemata (the building blocks), we have reduced the complexity of our problem; instead of building high-performance strings by trying every conceivable combination, we construct better and better strings from the best partial solutions of past samplings" (7:41).

Together, the schema theorem and the building block hypothesis help to explain why genetic algorithms work. However, the hypothesis does not give a formula for designing genetic algorithms. Instead, most practical information on genetic algorithm design and performance have come from empirical studies.

Permutation Genetic Algorithms

The simple genetic algorithm with a binary coding is appropriate for many unconstrained optimization problems. However, in many *order-based* problems, the solution may be specified by a specific arrangement of items. Examples include scheduling problems and the Traveling Salesman Problem (TSP).

Traveling Salesman Problem. The Traveling Salesman Problem is easily stated: a salesman must visit customers in each of n cities without visiting a city twice (13:165-167). The objective is simply to travel the least total distance and return to the starting city. A sample network of cities is shown in Figure 4. The solution to the TSP can be represented as a list of integer numbers with each integer corresponding to a city, and the cities visited in the order of the list. For example, a starting solution for a TSP with six cities could be represented as: $A = 1 \ 2 \ 3 \ 4 \ 5 \ 6$. This solution visits each city in ascending order from "1" to "6", returning to "1" to complete the tour. All possible solutions to the TSP can be

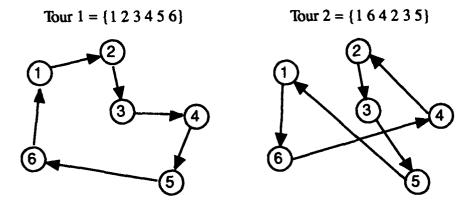


Figure 4. Sample TSP tours

represented as a permutation of the list of integers. A permutation is simply an arbitrary reordering of a set of items in a list. One permutation of the 6-city problem is A = 1 6 4 2 3 5. Although the ordering of the cities has changed, the number of cities remains the same and no city values are repeated in the list.

With n cities, there are $\frac{(n-1)!}{2}$ possible tours; a TSP with n=6 has $60 \left(\frac{5!}{2}\right)$ possible solutions. For a TSP twice as large (n=12), the total number of permutations is 239,500,800. The number of possible solutions is obviously incredibly large. Although smaller problems can be solved via deterministic graph search formulations, most large TSP problems are solved (non-optimally) using heuristic techniques which take advantage of distance information between the cities (12).

Blind Traveling Salesman Problem. Note that the TSP being solved by the genetic algorithm is harder to solve than the typical TSP because it does not use any distance

information to solve the problem. "In the blind traveling salesman problem, the salesman has the same objective with the added restriction that he is unaware of the distance he travels until he actually traverses a complete tour" (7:170). However, other ordering problems may not have "distance" information that can be exploited by a heuristic. As pointed out by Whitley, "This is important, because it means the method may be used on sequencing problems where there are no actual distances to measure, but rather only some overall evaluation of the total sequence" (19:137).

Since genetic algorithms had been applied successfully to other optimization problems, it seemed natural to attempt a genetic algorithm solution to the TSP. However, when standard genetic algorithms are applied to the TSP, they have difficulty in finding good feasible solutions (9). The major problem with a standard crossover operator is that most solutions are infeasible (i.e., 1 2 2 4 5 6) and the genetic algorithm probably converges on sub-optimal solutions (13:167). Some early attempts used a hybrid technique to combine genetic algorithm crossover with a repair mechanism to make valid tours (9). This approach was somewhat successful, but other approaches have been attempted using order-based crossover operators (13:168-191).

Order-based Crossover Operators. Another approach to solving ordering problems uses a different representation than the standard genetic algorithm. Instead of binary digits, an order-based crossover operator uses a chromosome that directly represents a solution. This chromosome is simply a list of integers; an example is a solution to the six-city TSP problem described previously (164235).

To exploit the information in the chromosomes, an order-based crossover operator, like simple crossover in the standard genetic algorithm, preserves part of the first parent while incorporating information from the second parent. The position of the genes is important. This contrasts with a standard genetic algorithm where their value is important. For the benefits of genetic algorithms to be realized for ordering problems, crossover operators must find other ways of combining information from two parents to build offspring with better fitness.

Although not strictly following the standard schema theorem for binary-based genetic algorithms, researchers have developed the concept of ordering schemata for order-based genetic algorithms (7:175-179). In an ordering problem, the absolute or relative positions of the items are important.

Ordering Schema. Ordering schema use a "don't care" symbol "!" to represent unfixed positions in a string. This is different from the "*" symbol presented earlier which represented unfixed values in a string. For example, in a position-based schema, the string (! ! 6 4 ! !) represents cities 6 and 4 in the third and fourth positions, with the "!" position filled arbitrarily from the remaining cities. Possible strings with this position schemata include: (5 1 6 4 2 3), (2 5 6 4 1 3), (3 2 6 4 5 1).

Another ordering schema uses relative ordering. As described by Davis(5:79):

It is important to understand that what is being passed back and forth here is not information of the form "node 3 is in the fifth position." Instead this operator combines information of the form "node 3 comes before node 2 and after node 7." The schemata in a n order-based representation can be written in just this way. Let us denote the nodes in a permutation by their indices. The chromosome (5 1 6 4 2 3) contains a number of schemata, including (5 4), (1 4 2 3), (5 6 3), and (5 1 6 4 2 3).

New crossover operators were developed that could work with ordering problems. Examples include Partially Matched Crossover (PMX) (7:154), position-based crossover (16:343), and edge-recombination (19). These methods attempt to retain the benefits of crossover while maintaining feasible solutions. Two of these operators are described below.

PMX. PMX was developed by Goldberg for use in solving TSPs. He describes the operator using a ten-city problem as an example (7:171). Each city is visited in ascending order; a sample permutation is: (1 2 3 4 5 6 7 8 9 10).

Under PMX, two strings (permutations and their associated alleles) are aligned, and two crossing sites are picked uniformly at random along the strings. These two points define a matching section that is used to effect a cross through position-by-position exchange operations.

To see this, consider two strings:

$$A = 984 \boxed{567} 13210$$

 $B = 871 \boxed{2310} 9546$

PMX proceeds by positionwise exchanges. First, mapping string B to string A, the 5 and the 2, the 3 and the 6, and the 10 and the 7 exchange places. Similarly mapping string A to string B, the 5 and the 2, the 6 and the 3, and the 7 and the 10 exchange places. Following PMX we are left with two offspring, A' and B':

$$A' = 98423101657$$

 $B' = 81015679243$

where each string contains ordering information partially determined by each of its parents (7:171).

Position-based Crossover. Another crossover operator, position-based crossover, was developed by Syswerda (16) and is called uniform order-based crossover by Davis (5). This operator attempts to preserve information about the relative ordering of elements in each of the parents. Several random positions are selected from each parent. These positions are inherited by one child. The other positions in this child are inherited in the order they appear in the other parent, skipping over all those included by the first parent. For example, using the strings used to illustrate PMX, first generate a bit string that is the same length as the parents:

Next, fill in some of the positions on Child 1 by copying them from Parent A wherever a "1" appears in the binary template.

A list of the elements in Parent A associated with "0" is made, and permuted to appear in the same order as in Parent 2. These permuted elements fill in the gaps in Child 1 to complete crossover.

List of elements associated with "0": (9,5,6,1,3,10)

Permuted in order of Parent 2: (1,3,10,9,5,6)

Child 1 = 18431079526

A similar process is used to create a second child, but with the roles of Parent A and B reversed.

Although PMX and position-based crossover are effective for ordering problems, they process different kinds of ordering schema. PMX tends to respect absolute city position, whereas position-based crossover tends to respect relative city position. Davis explains:

The information that is encoded here, however, is not a fixed value associated with a position on the chromosome. Rather, it is relative orderings of elements on the chromosome. Parent 1 may have a number of elements ordered relatively well. Uniform order-based crossover allows Parent 2 to tell Parent 1 that others of its elements should be ordered differently. The net effect of uniform order-based crossover is to combine the relative orderings of nodes on the two parent chromosomes in the two children (5:79).

These crossover operators are mentioned because they were used during the implementation of this research. They are effective for problems where ordering is important. PMX is more important where absolute position matters, while position-based crossover is more effective where a relative ordering is more important.

Genetic Algorithms in Scheduling

Many recent papers involving genetic algorithms deal with scheduling problems. Two main approaches to solving scheduling problems with genetic algorithms have been developed: direct chromosome representation and indirect chromosome representation.

Direct Chromosome Representation. Bruns advocates a direct chromosome representation for scheduling problems. In a direct problem representation, the production schedule itself is used as a chromosome. No decoding procedure is therefore necessary. The extended chromosome representation requires the construction of domain-specific recombination operators (4:355) and would therefore be problem-specific. In this case:

The sequence of the items within a chromosome is of no importance. To determine the quality of a chromosome any arbitrary evaluation function that has been used in traditional scheduling approaches is applicable without any prior transformation. (4:356)

According to Bruns, this approach should perform better than domain-independent operators, although he admits he cannot rely on any theory (4). It also involves creating custom crossover operators for each problem, which may be difficult. However, this approach has the advantage that the genetic algorithm is allowed to search the entire solution space, not just the ordering of the requests.

Indirect Representation. An indirect representation may also be referred to as a hybrid technique (7:202) where part of the problem is solved with the genetic algorithm and the other is solved using a deterministic routine. The first part is a sequencing problem, solved by the genetic algorithm, which orders each job request in a list. The second part is a schedule builder which takes each job, in the order of the list, and attempts to place each job request in a place in the schedule without overlapping another scheduled job. Obviously, those jobs earlier in the list are easier to schedule since more room in the schedule is available.

The schedule builder uses a simple rule to schedule each job; usually, a job is scheduled in the first position which meets the constraints of the problem. Although a schedule builder can use some information to find a good place in the schedule for each job, it is usually best to attempt the genetic algorithm search for good overall schedules by permuting the list of supports (16).

Much research on genetic algorithms in scheduling has been conducted recently in areas such as job shop scheduling and vehicle routing (3:452-459). For application to the satellite range scheduling problem, resource and sequence scheduling are more relevant. Two of these are reviewed by way of example in the next section.

Scheduling Examples

F-14 Test Range Scheduling. To solve a resource scheduling problem involving a test laboratory for F-14 fighter aircraft, Syswerda separated the genetic algorithm

from the specific problem (16:332). The list of items to be scheduled is represented as a string of numbers. The genetic algorithm permutes the order of the items in this string to find the best order in which to schedule the items. Then, given the ordered list of items produced by the genetic algorithm, a schedule builder program builds a feasible schedule. The schedule builder is merely a program which attempts to schedule each item in the order presented by the string. The number of items successfully scheduled is returned to the genetic algorithm as a fitness score. This type of approach to scheduling assumes that given a correct ordering of tasks, the schedule builder program can build the best schedule. Syswerda noted:

One thing that is clearly important, especially with regard to the greedy considerations of a single task, is the position of that task in the list. The closer the task is to the front of the list, the greater is its chance that it will be placed into the schedule ... if two tasks both require a scarce resource, the first task in the list may prevent the second from being scheduled, implying that the order of tasks is also important. (16:340)

Syswerda's implementation is interesting because he was able to satisfy many scheduling requirements such as priority of items and user preferences for scheduling days. Such flexibility is important for successful implementation of a scheduling solution.

Coors Scheduling. Much work has been done at Colorado State University on developing genetic algorithms for application to the Traveling Salesman Problem and some scheduling problems. Whitley developed a genetic algorithm called GENITOR (described in Chapter III, and used for this research), and Whitley and Starkweather developed an order-based crossover operator called genetic edge recombination for use in solving traveling salesman problems. Along with such theoretical developments, they also applied genetic algorithm solutions to a warehouse/shipping scheduler at Coors (15:74), and a production line scheduler at Hewlett-Packard (18:358–360).

From this research Whitley generalizes the application of genetic algorithms to scheduling problems in general:

... a broad class of scheduling problems can be viewed as sequencing problems. By optimizing the sequence of processes or events that are fed into a

simple schedule builder, optimization across the entire problem domain can be achieved. Schedules have not always been viewed this way because there has not existed a general purpose mechanism for optimizing sequences that only requires feedback about the performances of a sample sequence... Thus, a "genetic" approach to scheduling has the potential to produce some very general scheduling techniques, and could be the foundation of a general purpose approach to sequence scheduling. (18:358)

For this research, an approach similar to that used by Syswerda in the F-14 test range scheduling problem seems promising, as it does not require a custom chromosome. Hence, existing genetic algorithm packages, such as GENITOR, may be used for implementation. This approach only requires selecting an ordering crossover operator and constructing a schedule building program to simulate the scheduling operation. As described in the next chapter, such a strategy can be easily formulated.

Other Solution Efforts for Satellite Range Scheduling

Solutions to the satellite range scheduling problem have been studied recently. These include an effort by IBM, and two thesis efforts at the Air Force Institute of Technology (AFIT). These efforts used mixed-integer programming and heuristic approaches to find solutions.

Arbabi's Approach. The first study of automating satellite range scheduling took place during the 1981-84 time period when IBM conducted a study to determine the feasibility of automating satellite range scheduling (1:271-277). Arbabi concluded that a mixed-integer programming approach was not feasible for problems with more than 50 requests. Instead, he developed an approach called Continuous Time Scheduling (CTS). The procedure was not described in detail, as it is apparently proprietary, but it used a heuristic approach. This procedure reportedly scheduled 92% of the requests for one day (1:277).

Gooley's Approach. Gooley used both a mixed-integer programming approach and heuristic scheduling methods (8). The satellite range scheduling problem was successfully formulated as a mixed-integer program (MIP), but the number of integer variables

prohibited direct solution. To reduce the number of integer variables, Gooley divided the problem into two mixed-integer programs which scheduled low-altitude satellite supports (the MIP could handle up to 85 low-altitude supports at once). He then used heuristic insertion and interchange techniques to schedule the medium- and high-altitude satellite supports. Using this method, a test set of problems was solved in under 20 minutes with approximately 92% of all requested supports scheduled (8:5-2).

Schalck's Approach. In a follow-on thesis effort, Schalck improved on Gooley's solution by reducing the number of integer variables needed in the MIP formulation. By doing so, one MIP solution scheduled all low-altitude supports for one 24-hour period. Solution times for scheduling the high-altitude satellites in one 24-hour block was too long, and was reduced to about 30 minutes by scheduling high-altitude supports in two 12-hour blocks. The resulting solution scheduled approximately 98% of requested supports (this number is not directly comparable to Gooley's results because of differences in how the support requests were generated).

The mixed-integer programming approach taken by Gooley and Schalck is the best approach for the satellite range scheduling problem when the problem is small enough to be solved by the mixed-integer program, since such an approach should find an optimal solution. However, by decomposing the problem into separate problems, an overall optimal solution may no longer be guaranteed. For example, splitting the requests into two blocks means that support requests near the division point may not be scheduled. A genetic algorithm approach may produce better schedules while meeting the constraints of a short solution time because it would attempt to find a solution to the entire problem at once. Although the genetic algorithm approach is not guaranteed to find an *optimal* solution, it can attempt to find good solutions. Also, a genetic algorithm approach may be more flexible in handling additional constraints for special scheduling requests.

Summary

This chapter summarized the development of genetic algorithms including the standard genetic algorithm and variants. The extension to order-based genetic algorithms allows good solutions to some combinatorial optimization problems, such as the TSP. This in turn allows applications of genetic algorithms to scheduling problems by dividing the algorithm solution between a deterministic schedule builder and a genetic algorithm. Finally, past research efforts in satellite range scheduling were reviewed, with the conclusion that better solutions may be found by scheduling the entire day's schedule in one time block instead of decomposing the problem into smaller time blocks. This may be done with a genetic algorithm based approach.

III. Solution Methodology for Satellite Range Scheduling

This chapter presents a review of the satellite range scheduling (SRS) problem, followed by the development of a genetic algorithm-based scheduling strategy. By formulating the scheduling problem as a sequencing problem rather than as a mathematical program, a straightforward solution by an order-based genetic algorithm is possible. The chapter ends with a review of the implementation of the approach using the genetic algorithm package GENITOR.

Formulation

Definition of the Problem. In the SRS problem, satellite communication supports compete for RTS (remote tracking station) time in a 24 hour schedule. Each support must be scheduled in a restricted time window at certain RTSs due to visibility and scheduling requirements. For the low-altitude satellites, the requested support length fills the entire window. For medium to high altitude satellites, the time window is more flexible, as a tolerance for the beginning of the support is allowed for scheduling. Although each low-altitude support requires a fixed time window at one RTS and fills a time-window, the scheduling task is eased as most RTSs have two sides (antennas) capable of supporting communications. The RTS sides can support a satellite support simultaneously; if a satellite is visible to an RTS, it can be supported by any one of the available sides at the RTS. In addition, most medium-to-high altitude satellites are visible to more than one RTS. This fact, combined with the more flexible time windows of these satellites, makes scheduling them easier than scheduling low-altitude satellites.

Small Problem. An example set of time windows for a small set of five support requests is shown in Table 1. "Spt" is an arbitrary support number, "Begin" is the starting time for the window, "End" is the ending time for a window, "Length" is the actual service time needed, and "TAT" is the setup time required before a service time can begin. These supports are all serviced by RTS "POGO-A."

This sample time window data can be used to illustrate the satellite range scheduling (SRS) problem. This small problem is simplified since a real day's schedule would include

S ₇ :	Begin	End	Length	TAT
1	1	13	3	1
2	15	22	3	2
3	7	17	3	1
4	1	10	3	1
5	2	8	3	2

Table 1. Small Schedule Time Windows

over 300 supports, nearly all of which would have alternate windows for scheduling. These alternate windows could come from different antennas or sides at the same RTS, or from different RTSs. In addition to the table, this information is shown graphically in Figure 5.

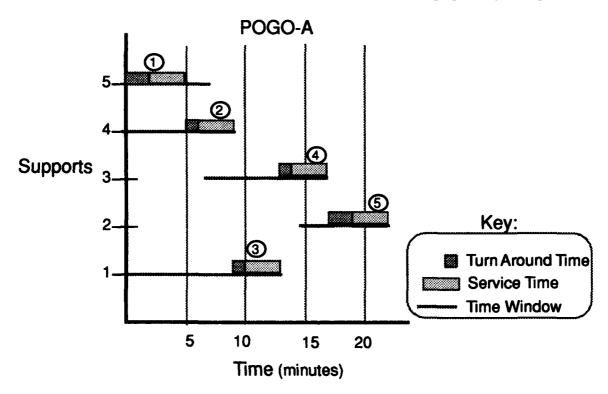


Figure 5. Small Schedule

The horizontal axis shows time in minutes, the vertical axis shows the supports by support number, and the entire chart is for RTS POGO-A. A problem with more than one RTS could be represented by more than one chart. In Figure 5, the *time window* is represented by a thin line, while the required TAT and service time is shown by the boxes (darker pattern for TAT). Such a representation clearly shows a schedule. The schedule

for the example successfully schedules all supports with no overlap of the support times and no violation of time windows. The TAT for support 5 is outside the service support window, but this is legal as no communications take place at that time, only set up for a support.

Mized-Integer Programming Approach. Previously, the SRS problem has been formulated as a mixed-integer programming (MIP) problem (Gooley and Schalck). For a problem with a small number of variables, this is the best approach because an optimal solution can be found. But, as the number of supports to be scheduled increases, so does the number of integer variables in the MIP formulation. Because of this increase, an optimal solution may no longer be available in a timely manner. Gooley and Schalck address this problem by decomposing the problem into parts. Gooley combined the use of a MIP formulation with heuristics. Schalck accomplished variable reduction and scheduled the satellites using an MIP for various combinations of satellites and blocks of time. These methods produce feasible schedules which are not necessarily optimal.

A MIP solution includes the starting time for each support and the RTS used for the support. The schedule can then be generated by adding the service times to the starting times. With this type of formulation, the mathematical program must find the actual starting times without violating constraints. A MIP solution of the small problem described in the previous section might begin with support 1 scheduled to start at time 10, support 2 scheduled to start at time 19, and so on, as shown in Figure 5.

Scheduling as a Sequencing Problem. Instead of formulating the satellite range scheduling problem as a mixed-integer program, this research approaches the problem as a sequencing problem where a sequence is defined as an ordered list of items. A solution to a sequencing problem in scheduling determines the "best" order in which to schedule items using simple rules for placing each item in the schedule. For example, in the satellite range scheduling problem, the solution is represented as a sequence of supports. In the small problem, there are five supports. An example ordering is: (5, 4, 1, 3, 2). To build a schedule from this representation, each support is scheduled according to the order it appears in the list.

In the example schedule, the first support, 5, would be placed in the first available time segment. The entire schedule is open and support 5 is scheduled at the beginning of its time window from time 2 to time 5. Note that the turn-around time of two minutes is scheduled from time 0 to time 2 which is outside the time window. This is valid since no communications take place during the turn-around time. The schedule is then updated to reflect that time 0 to 5 is no longer available. The second support in the list is 4. Its time window begins at 1, but it cannot be placed there because support 5 has already used minutes 0 to 5. Support 4 is then scheduled from time 6 to time 9 (1 minute TAT; 3 minutes for service). This support remains within its time window which ends at time 9. This procedure repeats for supports (1, 3, 2), and completes a schedule as shown in Table 2.

Spt	Begin	End	Length	TAT
5	2	5	3	2
4	6	9	3	1
1	10	13	3	1
3	14	17	3	1
2	19	22	3	2

Table 2. Small Problem Schedule

The actual start and end times for each support were determined by examining the completed schedule. Thus, all the information needed to define a solution is contained in the sequence of supports, and can be "decoded" by following simple deterministic rules. Obviously, such rules are somewhat arbitrary. The "first available" rule, as shown here, is simple. A rule which attempts to find the "best" place for each particular support would also work, but at the expense of simplicity and possibly time.

An arbitrary ordering of supports is not likely to produce a perfect, or even good, solution which schedules all support requests. However, if a solution exists, an ordering can be found which results in the optimal solution. In many cases, more than one ordering results in the same schedule depending on the interdependence of each support request. To find good schedules, an efficient way of generating better alternative orderings is needed. In this research, a genetic algorithm with order-based crossover operators is used to spawn

good orderings of the support requests. Genetic algorithms can quickly search the solution space, defined as the permutations of the list of items. These permutations can then be used to build a schedule as described above, and the number of supports scheduled successfully can be used as a fitness measure for the genetic algorithm search.

Hybrid Approach. As noted in the previous section, a schedule can be built from a sequence of supports using a schedule builder program. This schedule builder is required for building feasible schedules given a sequence of support by the genetic algorithm. The overall process is shown in Figure 6. Although some decision information can be incorporated into the schedule builder to improve local search, the genetic algorithm should do most of the exploration of the search space. This approach is based on the assumption that if the supports are entered in the schedule builder in a certain order, the greatest number of supports can be scheduled. It is reasonable to expect that many orderings of the supports may exist which produce the same schedule.

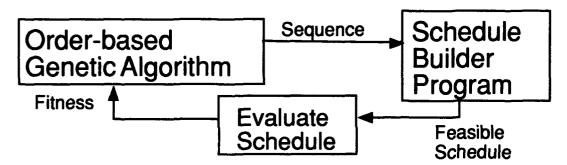


Figure 6. Schedule Builder Flowchart

The nature of the genetic algorithm and schedule builder approach allows other constraints to be imposed on the supports; for example, a request that a support be handled by a certain set of RTSs on a particular day, or that supports must be scheduled exactly one hour apart.

Algorithm Design and Implementation

Overview. Implementation of a hybrid approach requires data preprocessing of satellite requests to generate support request time windows, followed by application of the

GENITOR genetic algorithm. The schedule-building evaluation function is used within GENITOR to produce a given schedule, given a sequence of supports from a GA string.

Assumptions. The research assumes the constraints of the satellite range scheduling problem are those used in the previous efforts by Gooley and Schalck (14:1-7):

- 1. Requested support times are known in advance.
- 2. Time windows for satellite visibility are not flexible. The goal of the SRS problem is to schedule as many supports without having to change the time windows. After the initial schedule is formed, these constraints may be relaxed to schedule unscheduled supports. However, it is desired to minimize the number of such changes.
- Downtimes for RTS maintenance are not included. If known, flexible downtimes
 could be added as additional supports. Fixed downtimes could be be scheduled
 before regular supports are scheduled.

Use of these assumptions allows comparison of the genetic algorithm results to previous efforts.

Data Processing. Time windows for each support are processed from the raw satellite visibility and request data in the ASTROS database for a day. The support requests and time windows are represented as lines in a database with the following information:

Support Number: arbitrary support number

RTS: Remote tracking station and antenna side (example: POGO-A for A-side, POGO-B for B-side)

Beginning of Time Window: in minutes (example: 0200 is 120 minutes; 2400 is 1440 minutes)

End of Time Window: in minutes

Support Length (minutes): actual service time needed

Turnaround Time (TAT): set up time at RTS needed before service time (20 minutes for low-altitude satellites; 15 minutes for medium-high altitude satellites)

Satellite Identification: IRON (first four digits; identifies satellite) and revolution number (last three digits; identifies orbital pass)

An example set of time windows is shown in Table 3. This data is the same as for Table 1, but with fields included for the RTS name and and satellite identification number.

Spt	RTS	Begin	End	Length	TAT	Ident
1	POGO-A	1	13	3	1	2532097
2	POGO-A	15	22	3	2	4774042
3	POGO-A	7	17	3	1	9845009
4	POGO-A	1	10	3	1	3187074
5	POGO-A	2	8	3	2	9757024

Table 3. Small Schedule Time Windows

This time window data is read into data structures at the beginning of a run.

Data Structures. Time window data is read into a structure of arrays. The information can then be used each time a schedule is built from a sequence of supports. These array structures keep track of time window alternatives for each support, and allows the schedule to be filled in as supports are scheduled.

Schedule array. An array named filled (RTS, time) shows the status of each minute (time) for every RTS and antenna combination. A '0' represents an empty minute; a '1' indicates that minute is filled. This array is updated whenever a support is scheduled by setting the minutes used by the support to '1.' Any set up time is also blocked out of the schedule.

Support Information. Three arrays specify the requirements of each support: NumWin(support) is the number of time windows where a support can be potentially supported. Then the support time length requirements and turn around time are given by Length and TAT, respectively. This information is used when running the schedule builder program.

For each alternative time window for a support, three arrays give specific information.

BVIS and EVIS give the beginning time and ending time, respectively, of each time window.

RTS specifies the particular RTS used for the support.

Schedule Builder Program. The schedule builder is implemented as a program in the C language since many genetic algorithms are written in C. It is used as an evaluation function for the genetic algorithm by building a schedule from a sequence of supports generated by the GA, as shown in Figure 6. To construct the schedule builder, an insertion program was written, modeled loosely on those used by Gooley (8), which attempts to schedule each support into an available time window. Many alternative schedules are possible, as the insertion rules for the schedule builder are arbitrary. For example, the simplest implementation attempts to schedule supports to the first available numbered RTS. A schedule could also be built by assigning a support to an available RTS with the most open space left in its schedule. The "first available" approach was chosen for simplicity, and to schedule supports next to each other to utilize the available time at the RTSs. This approach simply schedules satellite communication supports to the first available position in time, and across each RTS, starting with the first RTS.

Each support is scheduled in the order it appears in the ordered list of support requests. The schedule builder tries the first RTS where there is a window for the support. If this fails, it tries the next until it runs out of windows. If the support is scheduled, the schedule score is incremented and the space used is blocked out of the schedule. Pseudocode for the program is as follows:

Empty schedule
Attempt to schedule each support
Try each window until succeed or exhaust windows
Try each set of empty spaces until succeed or reach end
Update schedule if support scheduled
Update score if support scheduled
Output final score

This can also be represented by a flow diagram, as in Figure 7.

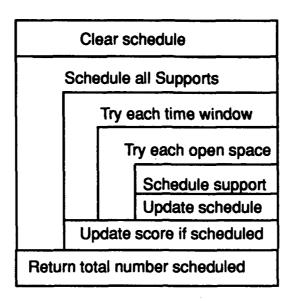


Figure 7. Schedule Builder program flow

This procedure finds the schedule corresponding to a sequence of supports. It is also used to build a schedule given a sequence of supports or to evaluate a given sequence as part of a genetic algorithm implementation.

Genetic Algorithm Implementation. To find "good" support sequences, the schedule builder is integrated into a genetic algorithm as an evaluation function. The genetic algorithm, GENITOR, and associated parameters are described,

GENITOR. Whitley's GENITOR (GENetic ImplemenTOR) code was chosen for use in this research because it has been used with success for other scheduling problems (17). GENITOR includes crossover operators for order-based genetic algorithms (edge recombination, order, and PMX) which have been shown to be useful in solving many scheduling problems (15). GENITOR differs from the standard genetic algorithm as described in Appendix A.

The evaluation function assigns a raw fitness measure to each population member. In GENITOR, the fitness is assigned by an evaluation function written in the C programming language (17). The function input parameter is the chromosome of a population member which is a list of supports. After attempting to schedule each support in the schedule

builder, the function returns a fitness value. In the satellite range scheduling problem, the chosen fitness measure is the number of supports successfully scheduled.

GENITOR Parameter Settings. In GENITOR, the major settings are the crossover operator, population size, selective pressure, and mutation rate (for binary-encoded chromosomes):

Crossover. The crossover operator is important in the quality of the answer.

GENITOR includes crossover operators for order-based genetic algorithms as discussed in Chapter II. These include edge recombination, order, and PMX.

Population Size. As explained in Chapter II, the population size affects the convergence rate of a genetic algorithm by controlling the variety of genes in the population. A smaller population may converge quickly, but usually to a sub-optimum. A larger population converges more slowly, and usually, but not always, finds a better final answer. For example, a population size of 30 finds a good answer within 1000 reproductions, but converges to a worse solution.

Selective Pressure. Selective pressure is a selection parameter specific to GENITOR. This parameter controls the rate at which a population converges by giving more reproductive opportunities to higher ranking individuals in the population. In GENITOR, selective pressure is usually set between 1.0 and 2.0. For example, a selective pressure of 1.5 gives the top-ranked individual 1.5 times the chance to reproduce than the median-ranked individual. A lower setting slows down convergence, hopefully allowing more time for the best solution to emerge. A higher setting drives the GA towards the final answer more quickly, but often at the expense of the best solution.

Mutation. It is important to keep a diverse genetic pool as the population converges. Mutation randomly changes part of a chromosome and is usually of less importance than crossover. Mutation was not used in this research because an order-based mutation operator was not present in GENITOR. The addition of such an operator may improve results and is discussed in Chapter V.

Day One Results. Various population sizes and selective pressure settings were tested in GENITOR using the first day of data to determine good parameter settings. In preliminary tests, all order-based crossover operators in GENITOR were tested. Position-based crossover found the best answer more quickly than the others. Position-based crossover was expected to produce good schedules since it explores the relative position of supports. Since supports are competing for resources, it makes sense that the relative order of supports in a scheduling list is more important than an absolute position in the list. Based on these results, the position operator was chosen as the operator for the test set, although the other order operators would be expected to perform nearly as well. The edge-recombination operator is tailored more towards the TSP and did not do as well as the other operators in this scheduling problem because it stresses adjacency information instead of relative ordering. This agreed with past results using edge-recombination for scheduling problems (15:74).

For Day One data, a graph of performance is shown in Figure 8. As GENITOR executes, at a given interval of reproductions it displays the best individual, the worst individual and the population average. The population has converged when these three numbers are equal, and little improvement can be expected. Convergence indicates that the individuals all have the same fitness score; they are either identical or similar enough to evaluate to the same fitness.

At any time during the run, the GA can be stopped and the best individual found thus far can be chosen as the best schedule found. Usually, the best schedule is found before convergence, but there is no way to know this in advance. In the Day One data, the best schedule was usually found at about 4000 reproductions, while the population did not converge until about 6100 reproductions. A genetic algorithm run could be stopped prior to convergence if time is critical.

Vary Parameters. The runs shown in Figures 9 and 10 display the interplay of exploitation versus exploration in a genetic algorithm solution (7:37). A high selective pressure or a small population size leads to quick convergence of the population. This exploitation of the best individuals found so far often leads to a quick, but less satisfactory

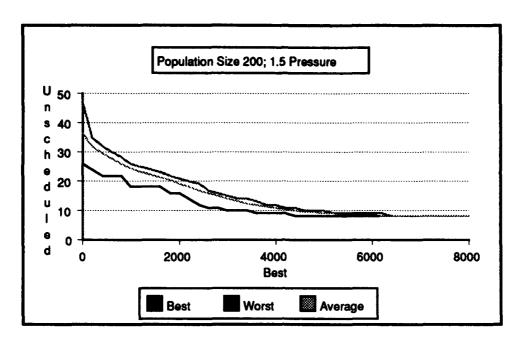


Figure 8. Day 1 Data Convergence

solution. On the other hand, a larger population with lower selective pressure encourages population diversity longer, but at the cost of slower convergence.

Repeated Runs. Since a GA has random components, a set of 10 runs was used for the Day One data in order to initially evaluate the variance of the results. Of 322 supports requested, the number of unscheduled supports was 8 in five of the runs and 9 in the other five runs; the average is 8.5, with a variance of 0.28. For this problem, the difference between runs is minor, and in practice one run should be sufficient for scheduling.

Random Schedule Generation. To give a baseline of the performance of this work, one might ask how well a random permutation of supports would perform. To answer this question, an additional run which generated random permutations of the support list and then used the schedule builder program to create schedules was performed. It found reasonable results within 8,000 reproductions, with a best solution of 23 supports unscheduled out of 322. This is not as effective as the best genetic algorithm schedule of only 8 unscheduled. These results are shown in Figure 11.

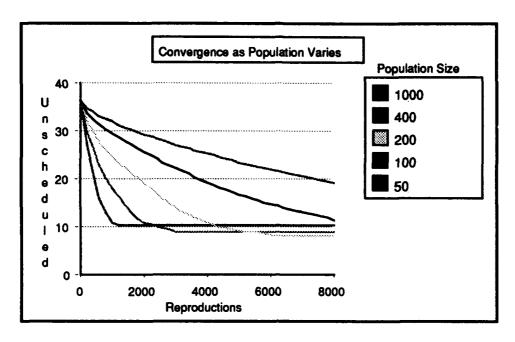


Figure 9. Vary Population Size for Day 1 Data

Overall Implementation. The overall implementation of the genetic algorithm-based approach begins with the processing of raw satellite request and visibility data into a format readable by the GENITOR main program. This processing is handled by Pascal programs written by Gooley and Schalck. After processing, the time window and request data for each support is stored in an array structure which the schedule builder uses to build feasible schedules.

Summary

This chapter describes how the SRS problem may be solved by scheduling items in an order determined by a genetic algorithm. By using this evaluation of each schedule to determine fitness in an order-based genetic algorithm, better schedules can be developed. The genetic algorithm package chosen for this research is GENITOR, because of its success in similar problems. Experiments using the first day of data provide good parameter settings for the GENITOR genetic algorithm in producing 24-hour schedules. The final settings are: position-based crossover, population size of 200, and selective pressure of 1.5. These settings gave a good compromise between the greatest number of supports scheduled and execution time. These settings are used to produce the schedules in Chapter IV.

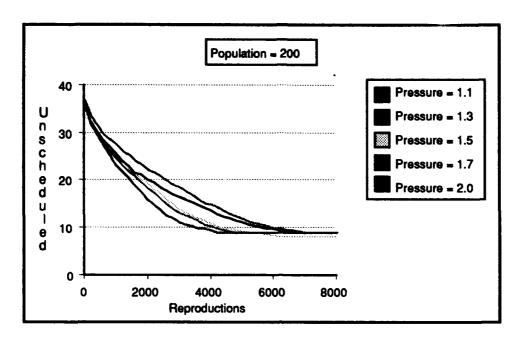


Figure 10. Vary Selective Pressure for Day 1 Data

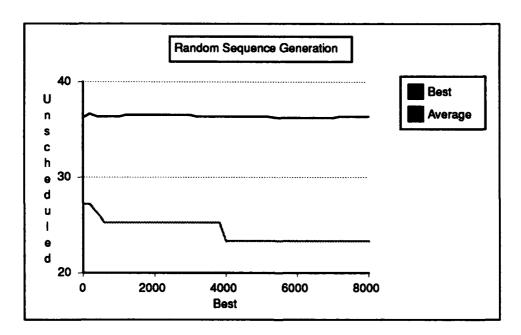


Figure 11. Random Sequence Generation for Day 1 Data

IV. Results

In order to evaluate the performance of a genetic algorithm approach to satellite range scheduling, the schedule builder model described in Chapter III was integrated into the GENITOR (17) genetic algorithm. This genetic algorithm was then tested on real satellite range data. In nearly all cases, the GA-based solution was able to match or exceed previous results that used a mixed integer programming (MIP) approach.

The primary purpose of this research was to evaluate the performance of the GA-based solution in scheduling satellite supports. A secondary purpose was to explore the flexibility of the GA-based approach in handling additional constraints on the scheduling of supports. Such flexibility would be important in implementing a genetic algorithm-based approach in the real world. The additional constraint tested here is the priority of different supports.

The criteria used to judge scheduling results is the number of supports scheduled and the time required for the solution. Supports not scheduled in the initial 24-hour schedule by the range schedulers must be deconflicted through coordination with the MCCs. In the worst case, those supports that cannot be rescheduled by relaxing constraints may have to be canceled. Short solution times are important if different parameters or priorities are to be used by the schedulers. Short solution times are also important for rescheduling on short notice.

Test Set

Scheduling data used in previous research efforts was available and used to test the hybrid GA approach, and to allow a comparison between these results and results of previous efforts. This data is taken from the ASTRO (Automated Scheduling Tools for Range Operations) (1:271) database which is currently used to assist manual scheduling. The primary data is from the seven day period from 12 Jul 92 to 18 Jul 92. This is the same data used by both Gooley and Schalck. Each day includes approximately 300 support requests and their associated request windows. Of these, approximately half are low-altitude satellite support requests.

The database contains satellite support requests and the visibility windows of the requests to remote tracking stations for each day. Pascal programs developed by Gooley and modified by Schalck were used to to put the requests in a format suitable for further work. Schalck's processing of the data was used here to allow comparisons to his results. After processing, the final data tables list the visibilities and requests for each day as shown in Table 3 of Chapter III. Processing details are available in Appendix D.

Although each day is scheduled separately, supports are allowed to overlap into the next day if their time window extends into the next day. This is included to match the approach of Schalck. The genetic algorithm implementation schedules across days by keeping track of the overlap into the next day and scheduling these overlaps in the next day's schedule before scheduling regular supports.

The schedule builder code developed in Chapter III serves as the evaluation function in GENITOR. The schedule builder produces a valid schedule from an ordered list of supports for each string in the population. The number of supports successfully scheduled is the fitness of each string. This fitness value is then returned to the genetic algorithm. The fitness is used by GENITOR to rank a particular schedule in relation to others in the population. The more fit strings have a higher probability of being selected to reproduce and exchange their genetic information with other strings through crossover.

Experimental Procedure

Tests were conducted in two major stages. The first day's data was tested to find good parameters for the GA (as shown in Chapter III). These parameters included population size and selective pressure. The primary goal of these parameters is to facilitate achievement of the best solution within the shortest time, where the solutions are scored by the number of supports scheduled. Good general settings were sought which would work well for the data sets. In practice, the general settings would not need to be changed each time a new set of data is introduced.

The final settings included position-based crossover, a population size of 200, and a selective pressure of 1.5. These settings appear to be robust in giving a good compromise between the best results and execution time.

These parameters were used to produce schedules for all seven days. Experimental runs were performed on Sun Sparc-10 and Sparc-2 workstations. Solution times for one day of data averaged 10 minutes on a Sparc-10 workstation and approximately 24 minutes on a Sparc-2 workstation. On the Sparc-10, producing 1,000 schedules took approximately one minute.

Week One Results. Once a good set of parameters was found, schedules for all seven days of data were produced. The genetic algorithm (GA) results are from GENITOR. with a population of 200, selective pressure of 1.5, and up to 8,000 reproductions. Overall results are shown in Figure 12 and Table 4. Supports whose time windows extend into the next day were allowed to be scheduled into the following day. These overlaps were then blocked out of the available RTS time for the next day. This is done by scheduling these overlap requests before the normal support requests in the following day's schedule. The corresponding results from Schalck's mixed-integer programming (MIP) solutions are included for comparison (14:4-2).

The total number of supports requested and scheduled are shown in Table 4. The difference between the MIP and GA-based results are labeled "GA-MIP."

Table	4. Re	sults:	for W	eek	One	Data	
ested	GA.	# Sch	edule	4	MTP-	# Sch	

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	322	314	312	2
2	302	296	296	0
3	311	307	304	3
4	318	315	311	4
5	305	301	299	2
6	299	292	292	0
7	297	291	291	0

The genetic algorithm results compare favorably to the MIP results of Schalck. In every case, the number of total supports the genetic algorithm schedules is at least as good as the number of supports scheduled by the MIP.

Breakdown by Support Type. Results are also broken down by low-altitude and high-altitude supports. As mentioned in Chapter I, low-altitude satellite supports are usually given higher priority in scheduling since they are more difficult to reschedule. Here the GA scheduled more high-altitude satellites than in the MIP solution and scheduled fewer low-altitude satellite supports than the MIP approach. This is due to the fitness of a schedule being calculated as the total number of supports scheduled without reference to the support type. Schalck's solution guaranteed scheduling of the greatest number of low-altitude supports possible because they were all scheduled before attempting to schedule any high-altitude satellites. This guarantees the scheduling of the greatest number of low-altitude satellites, but this may not always be necessary. Since set up times for low-altitude supports may be flexible, adjustment of the low-altitude satellite setup times may allow scheduling more supports overall.

Table 5. Low-altitude Results for Week One Data

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	153	147	149	-2
2	137	132	134	-2
3	146	143	143	0
4	142	139	140	-1
5	142	139	139	0
6	144	138	138	0
7	142	138	138	0

Additional Runs

Since one of the advantages of a hybrid GA approach is flexibility, other constraints on the scheduling of supports should be allowed. The first obvious constraint for inclusion is the scheduling of priorities for different support types.

Table 6. High-altitude Results for Week One Data

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	169	167	163	4
2	165	164	162	2
3	165	164	161	3
4	176	176	171	5
5	163	162	160	2
6	155	154	154	0
7	155	154	153	1

Support Priority. Low-altitude supports usually take precedence over high-altitude supports since they are more difficult to schedule. To introduce this preference in scheduling, different fitness scores can be assigned to the different types of supports. As an example, low-altitude supports were given a score of two for each one scheduled; the value of high-altitude supports was set to only one. Where a low-altitude support and a high-altitude support compete for the same position in the schedule, the low-altitude support should be scheduled. The overall results are shown in Table 7 and in Figure 13. A breakdown by low-altitude and high-altitude satellites is shown in Table 8 and Table 9.

Table 7. Week One Data with Low-Altitude Supports given Priority

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	322	313	312	1
2	302	296	296	0
3	311	307	304	3
4	318	314	311	3
5	305	300	299	1
6	299	293	292	1
7	297	291	291	0

Note that more low-altitude supports were accomplished, although at the cost of a tradeoff of some high-altitude supports. Note, however, for Day Two, the GA scheduled two fewer low-altitude supports than the MIP. Upon examination of the schedule, this is due to some high-altitude satellites from Day One scheduled into Day Two.

This type of priority assignment could also be applied to other situations; even so far as to give each satellite a different priority. An array of priority values could be used

Table 8. Low-altitude Results with Low-Altitude Supports given Priority

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	153	149	149	0
2	137	132	134	-2
3	146	143	143	0
4	142	139	140	-1
5	142	139	139	0
6	144	138	138	0
7	142	138	138	0

Table 9. High-Altitude Results with Low-Altitude Supports given Priority

Day	# Requested	GA# Scheduled	MIP# Scheduled	GA - MIP
1	169	164	163	1
2	165	164	162	2
3	165	164	161	3
4	176	175	171	4
5	163	161	160	1
6	155	155	154	1
7	155	153	153	0

to provide a different score for scheduling each support. A higher score would indicate greater scheduling priority for that support.

Two-Day Scheduling. Even with a priority score given to low-altitude satellite supports, the GA does worse in scheduling low-altitude supports for Days Two and Four. For Day 2, this is caused by the overlap from Day 1. Since each GA only schedules for one day, without regard to the following day, overlaps may reduce the overall number of supports scheduled. The simplest way to overcome this is to schedule two days at once in one block, with priority given to low-altitude supports.

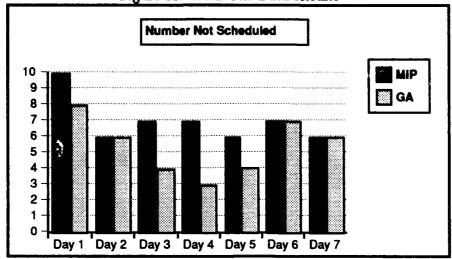
This was attempted with data from Day One and Day Two. Results were encouraging, as the overall number of supports scheduled remains the same, but the number of low-altitude satellite supports scheduled increased to the same level as the MIP solution. In return, the number of high-altitude supports scheduled decreased. Thus, the priority scheme, in conjunction with scheduling for two days, succeeded in matching the

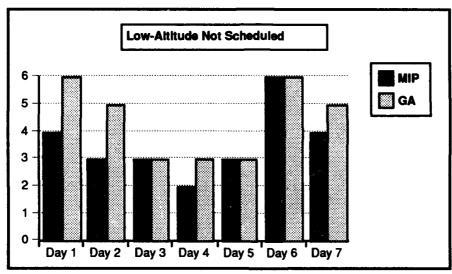
low-altitude results of the MIP. Note that the genetic algorithm in this case scheduled two days at once. The solution time was approximately double that for scheduling one day. This indicates that an increase in the number of supports for a day may be handled without much degradation in solution time or quality.

Summary

Results using one week of satellite support data indicates the GA can be successful in scheduling satellite supports. Results match or exceed those returned by the Schalck MIP approach. Solution times are short, and only one computer run is required.

Figure 12. Week One Data Results





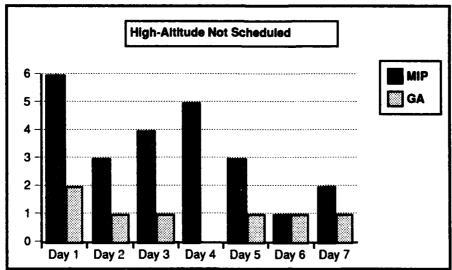
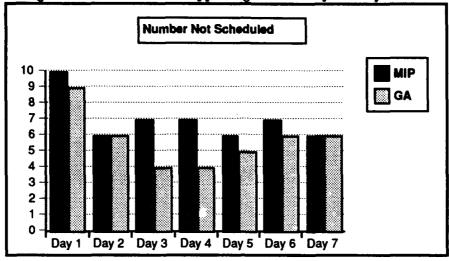
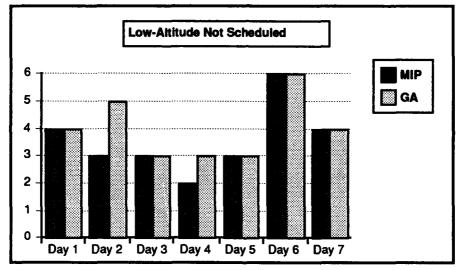
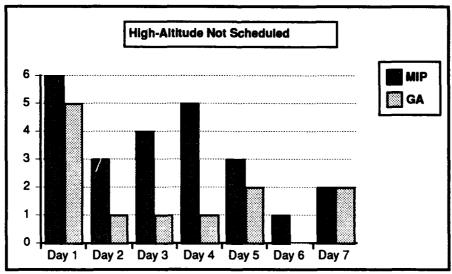


Figure 13. Low-Altitude Supports given Priority for Day 1 Data







V. Conclusions and Recommendations

Conclusions

The satellite range scheduling (SRS) problem involves scheduling over 300 satellite communication supports between Air Force Satellite Control Network (AFSCN) satellites and remote tracking stations (RTS) for a 24-hour period. Scheduling the greatest number of requests minimizes the time needed by schedulers to resolve conflicts in the schedule. The scheduling is currently done manually, with computer aid. Although the current system schedules approximately 95-98% of the requested supports, automating the scheduling process may produce better schedules in less time and use less human resources. The objective of this research was to discover if a genetic algorithm based approach can be effective in automating the scheduling of the requests for a 24-hour period.

The problem can be formulated as a mixed-integer program, but the size of realistic problems precludes producing optimal schedules in a timely manner. By partitioning the original mixed-integer problem into subproblems, Schalck and Gooley found good solutions (14:4-1), but since they do not solve the entire problem at once, there may be room for improvement.

For a test set of seven days of scheduling data, the GA solution matched and often exceeded the results of previous MIP efforts by solving the entire problem at once instead of decomposing the problem into subproblems. The algorithm scheduled over 96% of requested supports. Note, however, these results do not include RTS downtimes. Solution times on Sun Sparc workstations suggest a 24-hour schedule can be produced in under 15 minutes.

These results suggest a genetic algorithm (GA) based method can produce good schedules for the SRS problem. Although the genetic algorithm-based approach described in this research appears to be capable of producing good schedules in a short time, much work remains to be done to implement this automated scheduling approach. Further refinement of the program code may decrease time needed to provide a schedule. More importantly, a fully automated scheduler must also schedule other types of supports which occur in the scheduling process.

Recommendations

Improvements to the algorithm involve further modification and testing of the genetic algorithm with more complicated support requests. The genetic algorithm-based method is easy to implement and modify since all that is required is a schedule builder program to build a schedule one support at a time and to evaluate its fitness score. Since this schedule builder program is easy to test, modifications which allow new constraints can be tested quickly.

Minor Improvements.

Schedule Builder. If the GA-based algorithm is to be implemented, the program code for the schedule builder must be improved. Both readability and code efficiency would lead to better performance and easier modification by users. The program could also be integrated with the processing of satellite request data to make scheduling a one step process for the schedulers, from data processing to final schedule output.

Genetic Algorithm. The genetic algorithm implementation in GENITOR may be improved by adding a mutation operator. Although it does not play as large a role in the genetic search as crossover, mutation acts to keep variety in the population by randomly changing the genes of the population. An order-based mutation operator, such as that advocated by Davis (5:81), would be appropriate for use in this scheduling algorithm. Also, other genetic algorithm packages could be used besides GENITOR. The schedule builder program would simply be used as an evaluation function in the genetic algorithms.

Extensions. To make the produced schedules more realistic, other types of supports should be scheduled. RTS downtimes and scheduling priorities should be determined and included in the schedule. Also, the schedule builder program could be modified to schedule special requests not currently considered by this program. The most important of these would involve scheduling long requests by dividing them into shorter requests which may be scheduled on different RTSs.

Summary. A genetic algorithm based approach offers a way of quickly (minutes) scheduling 24-hour schedules in the satellite range scheduling process. In addition, the general approach could also be applied to other similar scheduling problems where resources must be scheduled. These include ICBM crew scheduling and test range scheduling.

Appendix A. Modified GENITOR Program Code

This appendix describes the GENITOR genetic algorithm, and lists the files modified to use GENITOR with the satellite range scheduling problem. These files include the main program and include files. Although not modified, the source code for the Position crossover operator is included since it is used for the final results.

Genetic Algorithm: GENITOR. The code from an existing genetic algorithm was used for this research. Whitleys GENITOR (GENetic ImplemenTOR) code was chosen because it has been used for similar problems (17). It includes crossover operators for order-based genetic algorithms (edge recombination, order, and PMX). These operators have been shown to be useful in solving many scheduling problems (15).

GENITOR differs from the simple genetic algorithm in two ways. First, instead of creating distinct generations, it creates only one new offspring at a time. Then, rather than replacing the parents with the offspring, GENITOR replaces the lowest ranking member. This approach has been called *steady-state* reproduction (5:35). It works by first selecting two parents from the population and producing two offspring. One offspring is randomly discarded. The other offspring replaces the lowest ranking member of the population. This new string is then ranked in relation to the fitness of the other members of the population, and inserted into the population. The new member can then compete for reproductive opportunities.

The second difference from the standard genetic algorithm is the way fitness is measured. Instead of relying on raw fitness measures, GENITOR uses relative ranking of population members to determine reproductive opportunities. Ranking prevents scaling problems associated with raw fitness values. Scaling problems occur when one individual is so much better that the others that it dominates the population too soon, causing permature convergence of the algorithm. To determine the reproductive opportunities of population members, GENITOR uses linear bias scaling. This linear bias is usually set between 1.0 and 2.0. For example, a bias of 1.5 means the top-ranked individual has 1.5 times the chance of the median-ranked population member to be selected to reproduce (15).

Modified Main program Code. The following listing is the main program code for GENITOR, modified for use with the satellite range scheduling problem. At the beginning of a run, this code reads in satellite time window data and overlap data. During execuation of the program, it calls tsp eval.c to evaluate schedules. Finally, it stores the best sequence in the file tourdata.

This main program is modified for use with the Position crossover operator.

```
/* MAIN Program - 31 Jan 94 */
  - 1 Feb, 31 Jan: read/write to files */
/* - 28 Jan: add in pre/post processing for overlap */
/*
  Copyright (c) 1990
/* Darrell L. Whitley
/* Computer Science Department
/*
   Colorado State University
/* Permission is hereby granted to copy all or any part of
/* this program for free distribution. The author's name
/* and this copyright notice must be included in any copy.
main_pos.c
 *******
               *********************
#include <stdio.h>
#include <ctype.h>
#include <string.h>
#include "ga_random.h"
#include "gene.h"
#include "ga_global.h"
#include "ga_params.h"
#include "ga_pool.h"
#include "ga_selection.h"
#include "ga_status.h"
#include "ga_signals.h"
#include "op_position.h"
 Declare evaluation function which tells you
 how "good" a particular gene's solution is.
 NOTE: the input parameters for the eval
      function should always be a gene and
  the gene's length
 #include "eval_pos.h"
int
main (argc, argv)
int argc;
char *argv[];
 int
           i,j;
 GENEPTR
          mom, dad, child;
 FILE
          *fp;
         **coord_array;
 int
 CITY_DATA *city_table;
             num_diffs; */
/* Character names for input/output */
```

```
int dayFum;
static char overlap[10] = {"overlap0"};
static char tourdata[10] = {"tourdata0"};
static char datafile[10] = {"datafile0"};
/*******
Setup signal handlers.
setup_signal();
Set the global parameters according to command line arguments.
argc--;
argv++;
parse_command_line (argc, argv);
Print Parameter Values
***************
fprintf (stdout, "\n");
print_params(stdout);
fprintf (stdout, "\n");
/*********
Input the day number
and concat to file names
printf("What is the day number (1-7)?");
scanf ("%u", &dayNum);
printf("Day Number is %u\n",dayNum);
overlap[7] += dayNum - 1;
tourdata[8] += dayNum;
                        /* Adds day number to filenames */
datafile[8] += dayNum;
printf("%s\n%s\n%s\n",tourdata,datafile,overlap);
/************
Seed the Random Number Generator
srandom(RandomSeed);
/***********************************
Allocate a genetic pool referenced by the global, Pool
if ( !(Pool = get_pool(PoolSize, StringLength)) )
fatal_error(NULL);
                     *******
Read in a description of the points to be toured and
create a representation of the distance between them.
 ************************************
/* make_dist_array (NodeFile, StringLength); */
/**************
 Assign schedule parameters to the program
 /* #include "data.h" */
Read in data
 *****************
 read_data (datafile,Pool->string_length); /* set = string length */
                                    /* 1/28: read overlap data */
 read_overlap_data (overlap);
/* for (i=1; i<=18; i++)
printf( "gRTS[%d]=%d, gBeginVis[%d]=%d, gEndVis[%d]=%d,
gReqLength[%d]=%d, gTurnAroundTime[%d]=%d \n",i, gRTS[i],i,gBeginVis[i],
 i,gEndVis[i], i, gReqLength[i],i,gTurnAroundTime[i] ); */
/* printf("Initialize genetic pool\n"); */
Initialize the genetic pool with data.
         init_pool (SeedPool, Pool, 0, Pool->size, tsp_eval);
```

```
/* printf("Sort Initial genetic pool\n"); */
Sort the initial genetic pool data.
************************
sort_pool (Pool);
/* printf("Allocate temporary storage for parents of reproduction\n"); */
Allocate temporary storage for parents of reproduction, and for child.
mom = get_gene (Pool->string_length);
dad
     = get_gene (Pool->string_length);
child = get_gene (Pool->string_length);
/* printf("Allocate a table to be used with the Order1 (Davis) Operator\n"); */
/*****************
Allocate a table to be used with the Order1 (Davis) Operator
 city_table = get_city_table (Pool->string_length);
/*******
Optimize !
*********/
for (/* CurrentGeneration already set :
    either intialized to 0 in its declaration OR initialized by a restart of a previous experiment */;
    CurrentGeneration < NumberTrials;
    CurrentGeneration++)
·
 Choose two genes for reproduction.
 *******************************
get_parents(mom, dad, Pool, linear, SelectionBias);
  /*********************
Call Order operator to create a child
***********
    /* printf("Start Order operator\n"); */
if (bitgen())
/* printf("After bitgen = 0\n"); */
position (mom->string, dad->string, child->string, Pool->string_length, city_table);
else
/* printf("After bitgen = 1\n"); */
position (dad->string, mom->string, child->string, Pool->string_length, city_table);
/* printf("So, kid, how good are you?\n"); */
     ******
So, kid, how good are you?
 ****************
    child->worth = tsp_eval (child->string, Pool->string_length);
/* printf("Insert new gene into population according to its worth\n"); */
Insert new gene into population according to its worth
 insert_gene (child, Pool);
   /***********************************
 If the StatusInterval parameter was set and this is the appropriate
time, print the population best, worst, mean, and average to stdout
                if (StatusInterval & !(CurrentGeneration % StatusInterval))
   show_progress (stdout, Pool, CurrentGeneration);
    **************************************
 If the DumpInterval parameter was set and this is the appropriate
 time, save the population and key parameters to disk for later
reference (or to restart execution later.
```

```
if (DumpInterval && !(CurrentGeneration % DumpInterval))
       dump_status(Pool, DumpBase);
/***********
 Summarize Results
  final_pool (FinalPool, Pool, CurrentGeneration);
  fprintf (stdout, "\n");
  printf("Right before print_pool\n");
  print_pool (stdout, Pool, 0, 1);
  printf ("After print_pool, before tourdata\n");
  /**********
 Print out tourdata to file
if (fp = fopen(tourdata, "w"))
              printf("Inside print tourdata\n");
  for (i=0;i < Pool->string_length; i++)
   fprintf(fp,"%d ",Pool->data[1].string[i]);
printf("Its past!\n");
              fclose(fp);
else
printf("Cannot print tourdata file");
 }
```

Ga global.h Include files. This file contains the global variables used in GENITOR. The global variables used to store time window information for each support are included here as global variables and used in the tsp eval function. Another related file called ga global external.h is also listed.

```
/*ga_global.h */
/* mod 9 Feb: add temp LF Num for Priority */
/* modified 8 Feb for downtimes; 28 Jan for overlap */
/* Copyright (c) 1990
                                        */
/* Darrell L. Whitley
                                        */
/* Computer Science Department
                                        */
/* Colorado State University
                                        */
/* Permission is hereby granted to copy all or any part of */
/* this program for free distribution. The authors name */
/*
  and this copyright notice must be included in any copy.
                                        */
* These REQUIRED and OPTIONAL parameters can be chosen by the user. *
/* REQUIRED */
int
     PoolSize;
```

```
int
        StringLongth;
 int
        NumberTrials;
char
        WodeFile[80];
                           /* contains coordinates of tsp nodes */
/* OPTIONAL */
long Random
        RandomSeed
float
        SelectionBias:
        MutateRate = 0.0;
 float
 int
        StatusInterval = 0; /* dump status every nth generation */
 int
        DumpInterval = 0;
                           /* save state of every nth generation */
char
                           /* file containing init data */
        SeedPool[80];
        FinalPool[80];
char
                           /* file containing final results */
 char
        DumpBase[80];
                           /* basename of file(s) into
                              which to dump population */
        NumPop;
 int
                           /* Number of Subpopulations */
 int
        SwapInterval;
                           /* Trials between Swapping of Subpopulations */
 int
        SwapNumber;
                           /* Number of Strings Swapped between Subpops */
 int
                           /* Experiments must be used in main() */
        Experiments;
float
                           /* Cutoff value for a given experiment */
        CutOff;
int SequenceFlag = 0;
                           /* if set to 1 in main, insert_unique_gene() */
    /* will use different criteria for sameness */
        CurrentGeneration = 0;
POOLPTR Pool;
/* Schedule Builder Global Variables */
/*
    [spt][win] */
       gRTS[700][20];
int
       gBeginVis[700][20];
int
int
       gEndVis[700][20];
int
       gReqLength[700];
       gTurnAroundTime[700]:
int
int
       gBeginSched[700];
int
       gEndSched[700];
int
       gWumWin[700];
int
       gOverlapBegin[20];
int
       gOverlapEnd[20];
int gDownNum;
int gDownRTS[100];
int gDownBegin[100];
int gDownEnd[100];
int gNumLF;
/* ga_global_extern.h */
/* mod 9 Feb: add temp NumLF for priority */
/* mod 8 Feb for downtime */
/* modified 28 Jan 94 to add overlap variables */
/*
/* Copyright (c) 1990
                                                          */
/* Darrell L. Whitley
                                                          */
/* Computer Science Department
                                                          */
/* Colorado State University
                                                          */
/* Permission is hereby granted to copy all or any part of */
  this program for free distribution. The authors name
/*
   and this copyright notice must be included in any copy.
/*
```

```
extern int
              PoolSize;
extern int
              StringLength;
extern int
              NumberTrials;
extern char
              WodeFile[];
extern long
              RandomSeed
extern float extern float
              SelectionBias:
              MutateRate;
              StatusInterval;
extern int
extern int
              DumpInterval;
extern char
              DumpBase[];
              SeedPool[];
extern char
extern char
              FinalPool[];
extern int extern int
              CurrentGeneration:
              MumPop;
              SwapInterval;
extern int
              SwapHumber;
extern int
 extern int
              Experiments;
 extern float
              CutOff;
              SequenceFlag;
 extern int
extern POOLPTR Pool;
/* Schedule Builder Global Variables */
              gRTS[700][20];
extern int
             gBeginVis[700][20];
extern int
             gEndVis[700][20];
extern int
             gReqLength[700];
extern int
extern int
             gTurnAroundTime[700];
extern int
             gBeginSched[700];
              gEndSched[700];
extern int
extern int extern int
              gNumWin[700];
              gOverlapBegin[20];
extern int
              gOverlapEnd[20];
extern int gDownNum
extern int gDownRTS[100];
extern int gDownBegin[100];
extern int gDownEnd[100];
extern int gFumLF;
```

Position Crossover Operator

This program code is unchanged, but included since it was used in GENITOR to test the data.

```
**********************
 * This is the position operator developed by Syswerda *
   (The Genetic Algorithms Handbook, L Davis, ed) and *
   implemented by Susan McDaniel at Colorado State *
    University.
 * To use this program include the following lines in *
 * the main_tsp.c file: *
   This call should only happen once: *
   get_city_table (Pool->string_length); *
 * This code should be in a loop so it is called once *
 * for each recombination: *
* if ( bitgen () ) *
      position (mom->string, dad->string, child->string, *
                 Pool->string_length, city_table); *
      position (dad->string, mom->string, child->string, *
      Pool->string_length, city_table); *
 * This operator attempts to preserve position *
 * information during the recombination process. *
* Several random locations in the tour are selected
* along with one of the parents and the cities in *
 * those positions are inherited from that parent.
 * The remaining cities are inherited in the order in *
 * which they appear in the unselected parent skipping *
 * over all cities which have already been included in *
 * the offspring. *
 * Example- Position based crossover: *
 * Parent 1: a b c d e f g h i j *
 * Cross Pts:
                                    (Parent 1 selected) *
 * Parent 2: cfajhdigbe *
 * Offspring: a b c j h f d g i e *
 * The cities b, c , f and i are inherited from * * Parent 1 in positions 2, 3, 6 and 9 respectively.
 * The remaining cities are inherited from Parent 2 as * follows: Off[1] = P2[3] since P2[1] and P2[2] have *
 * already been included in the child tour. Then
 * going through Parent 2 in order, Off[4] = P2[4], *
* Off[5] = P2[5], Off[7] = P2[6], Off[8] = P2[8] and *
 * Off[10] = P2[10]. *
                    #include <stdio.h>
#include <malloc.h>
#include "ga_random.h"
#include "gene.h"
#include "op_position.h"
* FUNCTION: get_city_table
 * DESCRIPTION: allocates space for the city data table
 * the city data table contains the position
 * of each city in each parent and also
 * has a field which is set when a city has
```

```
* been included in the offspring tour.
 * IMPUT PARAMETERS: string length
 * RETURN VALUE: the address of the city table
 CITY_DATA *
get_city_table (length)
int length;
   CITY_DATA +city_table;
/* malloc one extra location so that cities 1-N can be accessed directly. location 0 will not be used \*+/
   if (!(city_table = (CITY_DATA *) malloc ((length + 1)*sizeof (CITY_DATA))))
      printf ("get_city_table: Malloc failure. \n");
   return (city_table);
* FUNCTION: position
 * DESCRIPTION: performs position crossover operation
 * IMPUT PARAMETERS: two parent gene strings
       space for child string
       the length of the two strings
       the city table address
 * RETURN VALUE:
 void position (dad, mom, kid, length, city_table)
GENE_DATA dad [], mom [], kid [];
int length;
CITY_DATA *city_table;
   int num_positions;
   int i, pos, dad_index, kid_index;
   for (i=1; i<=length; i++) { /* initialize city table */
      city_table[i].used = 0;
    /* select *positions that will be inherited directly from parent */
   num_positions = randomain (2*length/3, length/3);
   for (i=0; i<num_positions; i++) {</pre>
      pos = randomain (length - 1, 0); /* select position randomly */
      kid[pos] = mom[pos]; /* transfer cities to child */
      city_table[mom[pos]].used = 1; /* mark city used */
   dad_index = 0;
kid_index = 0;
   while (kid_index < length) {
      if (!city_table[mom[kid_index]].used) { /* next position in kid filled*/
  if (!city_table[dad[dad_index]].used) { /*next city in dad not used*/
    kid[kid_index] = dad[dad_index]; /* inherit from dad */
    dad_index ++; /* increment indexes */
    kid_index ++;
```

```
else { /* next city in dad has been used */
    dad_index ++; /* so increment dad index */
} /* end else */
} /* end if */
else { /* next position in kid is filled */
kid_index ++; /* so increment kid index */
} /* end else */
} /* end while */
}
```

Appendix B. Evaluation Function Code

The tsp eval program is a procedure which builds and evaluates a schedule. It takes as input an ordered list of supports to schedule and then, in the order of the list, attempts to schedule each support into a schedule. The function returns a fitness value to the main genetic algorithm.

The schedule builder code includes three procedures: ScheduleWindow, Schedule-Support, and tsp eval. Tsp eval is the main procedure. It is called from GENITOR, and returns the evaluation of a schedule. Tsp eval first clears a schedule, schedules any overlaps from the previous day, then attempts to schedule each support in the order given in the sequence passed to it from GENITOR. ScheduleSupport attempts to schedule each support in the available time windows. Each time window is attempted by calling the ScheduleWindow routine. If a support is scheduled in a window, the space taken by the support is blocked out of the schedule. Then ScheduleSupport returns a fitness value to tsp eval. This fitness value is normally one, but can differ if a priority scheme is used to weight the scheduling of different types of supports. After all supports have been scheduled (or failed to be scheduled), tsp eval returns a fitness score based on the number of supports scheduled to the main GENITOR function.

```
/* SCHEDULE BUILDER PROGRAM */
/* 31 Jan 94: auto file generation */
/* 28 Jan 94 Eval_tsp.c -- position crossover */
/* 28 Jan: add overlap variables */
/* This modification allows supports to be scheduled past the end of the day */
/* if allowed by the supports windows */
/* do this by increasing filled array from 1470 to 2100, and take off end */
/* conditions
#include <stdio.h>
#include <math.h>
#include "gene.h"
#include "op_position.h"
#include "ga_global_extern.h"
int
ScheduleWindow (spt, win, filled)
int
      spt;
int
      win;
int filled[21][2100];
     /* Declare Variables */
     int
          bVis, /* section /* section /* ending visibility */
                           /* gBeginVis - gTurnAroundTime : 0 if latter <0 */</pre>
    eVis,
                        /* RegLength + TurnAroundTime */
          rLength.
         endSpt,
                     /* Flag for Support either scheduled or run out */
```

```
/* Counter for support index being scheduled */
        j.
1,
                          /* Counter for filling in scheduled blocks */
                                     Index for support */
                    spt:
                          /* current time interval */
         minute,
                     /* length so far */
        length,
         scheduled,
                   /* binary: 1 or scheduled; 0 for not schedule */
         beginSched,
         endSched,
         start,
                     /* current start of support */
tat,
             /* turn around time */
                     /* RTS number -- passed from global gRTS[support] */
         rts;
                       /* scheduledPtr : indicates if scheduled or not */
     /**** Initialize Variables *****/
     scheduled = 0;
    length = 0;
     endSpt =0; /* initialize flag; end when it = 1*/
    /* Initialize Starting Point, Required Length with TAT */
    rts = gRTS[spt][win];
    tat = gTurnAroundTime[spt];
    bVis = gBeginVis[spt][win] - gTurnAroundTime[spt] + 20 ;
     eVis = gEndVis[spt][win] + 20;
/* add twenty for overlap area */
     if (bVis < 0)
{
           tat = bVis + tat;
  bVis = 0;
                  /**** reduces tat if near boundary at start ****/
}
    rLength = gReqLength[spt] + tat;
    minute= bVis + 1;
                       /* my blocks are filled to the right of position */
     start = minute;
                                 /* Starting position of current attempt */
     while (endSpt != 1)
                              /* schedule until successful or end */
          if (filled[rts][minute] == 0)
                                           /* if current space free */
                                         /* increment length if space free */
               length += 1;
          èlse
               length = 0; /* reset to zero if space is filled */
     tat = gTurnAroundTime[spt];
         rlength = gReqLength[spt] + tat;
               start = minute+1;
          if (length == rLength)
                                  /* then schedule the support here */
                   beginSched = minute - gReqLength[spt] - 20;
   endSched = minute - 20;
                   scheduled = 1;
                   for (l = start;l <= minute;l++)</pre>
                      filled[rts][1] = 1;
  }
          if ( (minute == eVis) || (scheduled == 1) )
             endSpt = 1;
          minute += 1;
     } /* End Support */
      /* fprintf (stdout, "scheduled = %d\n", scheduled); */
      /*if(scheduled == 1)
        printf("SCHEDULE Support %d in window %d at RTS %d, begin: %d
   end: %d\n",
                    spt, win, rts, beginSched, endSched );
     else
       printf("
                     Support %d not scheduled in window %d\n", spt, win); */
```

return(scheduled);
} /* end ScheduleWindow */

```
int
ScheduleSupport(spt,filled)
int spt;
int filled[21][2100];
int i, win, flag;
            /* flag = 1 if scheduled , 0 ow */
int sched;
sched = 0:
                   /* counter for number of windows tried */
win = 1;
flag = 0;
while (flag != 1)
     sched = ScheduleWindow(spt,win,filled);
    /* printf ("sched=%d, gWumWin[%d]=%d\n", sched,spt,gWumWin[spt]); */
if ((sched == 1) || (win == gWumWin[spt]))
flag =1;
    win +=1;
/* if (sched==0)
 printf("*** Support %d NOT SCHEDULED!***\n",spt); */
return(sched);
float
tsp_eval (order, num_supports)
GENE_DATA order[];
         num_supports;
int
{
     /* Declare Variables */
        filled[21][2100].
                           /* time units in schedule are filled[rts][time] */
                         /* Counter for support index being scheduled */
        Ř,
                       /* Index for support */
        spt,
        numScheduled,
                         /* number of supports scheduled */
                          /* counter for current time */
         minute,
        rtsNum;
                         /* max number of RTSs */
     /**** Initialize Variables *****/
    rtsHum = 20; /* must be one less than max cause of arrays 0-19 */
     /* printf("rtsNum = %d\n", rtsNum); */
    for (k = 1;k <= rtsNum;k++)
        for (minute = 0;minute <= 2100;minute++)</pre>
            filled[k][minute] = 0;
           /* printf("filled[%d][%d] = %d\n",k,minute,filled[k][minute]); */
        }
     /* Schedule the overlaps first, before the real supports */
     for (k = 1; k <= rtsHum; k++)
 for (minute = gOverlapBegin[k]; minute <= gOverlapEnd[k]; minute++)</pre>
     filled[k][minute] = 1;
     numScheduled = 0;
     /* Schedule each Support in Priority Order */
                                      /* Try to schedule support */
     for (j = 0;j< num_supports; j++)</pre>
          spt = order[j];
```

```
/* printf( "Start SS Order[%d] = %d\n",j, order[j]); */
         numScheduled += ScheduleSupport(spt, filled);
          /* end num_supports */
   /* fprintf (stdout, "num_supports = %d\n", num_supports); */
  /* The GA wants to MINimize something... so do Max poss - actual */
     /* printf( "\n Number Scheduled = %d\n****************\n\n\n",
numScheduled); */
    return(num_supports - numScheduled);
     /* End Eval */
FUNCTION: read_data
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support data
* RETURN VALUE: none
* CALLS: get_coord_array
         read_coords
get_2D_dist_array
         calc_distances
*****
             void
read_data (coord_file, num_supports)
char
              coord_file[];
int
             num_supports;
 int i;
 int spt;
 int rts;
  int bVis, eVis, rLength, tat;
 FILE *fp:
 printf("coord_file is: %s", coord_file);
 for (i=1; i<=num_supports; i++)</pre>
gWumWin[i] = 0;
                   /* set number of windows [spt] = 0 */
/* printf("gWumWin[%d]=%d\n",i,gWumWin[i]); */
 if (fp = fopen(coord_file, "r"))
        while (fscanf(fp, "%d %d %d %d %d %d", &spt, &rts, &bVis,
&eVis, &rLength, &tat ) !=EOF)
             gWumWin[spt] +=1;
             gRTS[spt][gNumWin[spt]] = rts;
             gBeginVis[spt][gMumWin[spt]] = bVis; /* add for overlap */
             gEndVis[spt][gMumWin[spt]] = eVis; /* add for overlap */
             gReqLength[spt] = rLength;
             gTurnAroundTime[spt] = tat;
             /* printf ("spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
 TAT = %d\n", spt, gRTS[spt][gNumWin[spt]], gBeginVis[spt][gNumWin[spt]],
gEndVis[spt][gWumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
     /* printf("----gNumWin[%d] = %d\n", spt,gNumWin[spt]); */
 fclose (fp);
 /* for (i=1; i<= num_supports;i++)</pre>
printf("gWumWin[%d] = %d\n",i,gWumWin[i]); */
 else
```

```
fatal_error ("Cannot read input city file datafile");
FUNCTION: read_overlap_data
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support overlap data
 RETURN VALUE: none
 * CALLS: get_coord_array
         read_coords
get_2D_dist_array
         calc_distances
                        *************
void
read_overlap_data (coord_file)
               coord_file[];
 int i;
int rts;
int bVis, eVis;
 FILE *fp;
 if (fp = fopen(coord_file, "r"))
        while (fscanf(fp, "%d %d %d", &rts, &bVis, &eVis ) !=EOF)
             gOverlapBegin[rts] = bVis;
     gOverlapEnd[rts] = eVis;
               /* printf("gOverlapBegin[%d] = %d\t gOverlapEnd[%d] = %d\n", rts,
 gOverlapBegin[rts], rts, gOverlapEnd[rts]); */
             /* printf ("spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
TAT = %d\n", spt, gRTS[spt][gWumWin[spt]], gBeginVis[spt][gWumWin[spt]], gEndVis[spt][gWumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
      /* printf("----gWumWin[%d] = %d\n", spt,gWumWin[spt]); */
 fclose (fp);
  /* for (i=1; i<= num_supports;i++)</pre>
printf("gWumWin[%d] = %d\n",i,gWumWin[i]); */
 else
fatal_error ("Cannot read input city file overlap");
```

Appendix C. Schedule Builder

This program code is a modification of the Evaluation Function code in Appendix B. Schedule Builder generates a schedule corresponding to an input solution sequence. This solution is stored in a file named *tourdata* and is produced by the GENITOR code at the end of a run.

The output of the Schedule Builder code is an overlap file and a schedule file. The overlap file contains the start and end times of all supports which overlap into the next day or within the maximum turn around time of the end of the day. The schedule file contains the schedule produced for the 24-hour period. Included in this file is the support number, RTS, beginning time, ending time, service time, and turnaround time.

```
/* Schedule Builder Code - build schedule */
/* 1 Feb: print schedule to file */
/* 30 Jan 94: fix overlap by clearing the global overlaps */
/* 28 Jan 94: overlap of 20 minutes allowed at beginning */
/* 3 Jan 94: allowed supports to be scheduled past 1470 */
/* changed filled[] array from 1470 to 2100, take off condition */
#include <stdio.h>
#include <math.h>
#include <gene.h>
#include <op_edge_recomb.h>
#include <ga_global_extern.h>
ScheduleWindow (spt, win, filled)
int
      spt;
int
      win;
int filled[21][2100];
     /* Declare Variables */
     int
                                         gTurnAroundTime : 0 if latter <0 */
          bVis.
                           /* gBeginV_
                           /* ending visibility 28 Jan */
          eVis,
                        /* RegLength + TurnAroundTime */
         rLength,
         endSpt,
                     /* Flag for Support either scheduled or run out */
         j,
                          /* Counter for support index being scheduled */
                          /* Counter for filling in scheduled blocks */
                    spt:
                                     Index for support */
         minute,
                          /* current time interval */
                     /* length so far */
         length,
                    /* binary: 1 or scheduled; 0 for not schedule */
         scheduled,
         beginSched,
         endSched,
                    /* current start of support */
         start,
 tat,
             /* turn around time */
                     /* RTS number -- passed from global gRTS[support] */
         rts;
                       /* scheduledPtr : indicates if scheduled or not */
FILE *fp;
     /**** Initialize Variables *****/
     scheduled = 0;
     length = 0;
```

```
endSpt =0; /* initialize flag; end when it = 1*/
      /* Open file for output to schedule 1 Feb */
fp = fopen( schedule ,
     /* Initialize Starting Point, Required Length with TAT */
    rts = gRTS[spt][win];
     tat = gTurnAroundTime[spt];
     bVis = gBeginVis[spt][win] - gTurnAroundTime[spt] + 20;
     eVis = gEndVis[spt][win] + 20; /* 20 is max TAT - 28 Jan */
/* if ((spt==1) || (spt==155))
printf( tat=%d,bVis=%d\n ,tat,bVis); */
     if (bVis < 0)
           tat = bVis + tat;
  bVis = 0;
                   /**** reduces tat if near boundary at start ****/
    rLength = gReqLength[spt] + tat;
/* if ((spt==1) | (spt==155))
printf( rLength=%d\n ,rLength); */
    minute= bVis + 1; /* my blocks are filled to the right of position */
     start = minute;
                                /* Starting position of current attempt */
     while (endSpt != 1)
                            /* schedule until successful or end */
          if (filled[rts][minute] == 0)
                                          /* if current space free */
              length += 1;
                                         /* increment length if space free */
          else
               length = 0; /* reset to zero if space is filled */
  tat = gTurnAroundTime[spt];
         rLength = gReqLength[spt] + tat;
               start = minute+1;
          if (length == rLength)
                  beginSched = minute - gReqLength[spt] - 20;
           endSched = minute - 20;
                   scheduled = 1;
                   for (1 = start; l <= minute; l++)</pre>
                      filled[rts][l] = 1;
                   /* If ends in TAT zone them save to overlap */
                                    /* 1440 - max TAT; in overlap */
           if (endSched > 1420)
              if ( (gOverlapEnd[rts] = 0) || (endSched >
 gOverlapEnd[rts]) )
        { /* if rts overlap is new or this one is longer than
 last */
     gOverlapEnd[rts] = endSched - 1420;
     if (beginSched > 1420 )
          gOverlapBegin[rts] = beginSched - 1420;
       gOverlapBegin[rts] = 0;
}
          if ( (minute== eVis) || (scheduled == 1) )
             endSpt = 1;
          minute += 1;
     } /* End Support */
      /* fprintf (stdout, scheduled = %d\n , scheduled); */
      if(scheduled == 1)
{
         printf( %d %d %d %d %d %d\n
                    spt, rts, beginSched, endSched , gReqLength[spt],
```

```
gTurnAroundTime[spt]);
fprintf(fp, %d %d %d %d %d %d\n
                   spt, rts, beginSched, endSched, gReqLength[spt],
spt,
gTurnAroundTime[spt] );
}
     /*else
       printf(
                     Support %d not scheduled in window %d\n , spt, win); */
    fclose(fp); /* stop printing to file 1 Feb */
    return(scheduled);
} /* end ScheduleWindow */
/*************
                         ************************
int
ScheduleSupport(spt,filled)
int spt;
int filled[21][2100];
int i, win, flag;
            /* flag = 1 if scheduled , 0 ow */
int sched;
sched = 0;
win = 1:
                   /* counter for number of windows tried */
flag = 0;
while (flag != 1)
     sched = ScheduleWindow(spt,win,filled);
                                           , sched, spt, gWumWin[spt]); */
    /* printf ( sched=%d, gHumWin[%d]=%d\n
    if ( (sched == 1) || (win == gNumWin[spt]) )
flag =1;
    win +=1;
/* if (sched==0)
 printf( *** Support %d NOT SCHEDULED!***\n ,spt); */
return(sched);
tsp_eval (order, num_supports)
int order[];
int
         num_supports;
{
     /* Declare Variables */
        filled[21][2100],
                          /* time units in schedule are filled[rts][time] */
                         /* Counter for support index being scheduled */
         j,
                       /* Index for support */
        spt,
        numScheduled,
                         /* number of supports scheduled */
         minute,
                          /* counter for current time */
        rtsNum;
                         /* max number of RTSs */
     /**** Initialize Variables *****/
    rtsNum = 20; /* must be one less than max cause of arrays 0-19 */
     /* printf( rtsNum = %d\n , rtsNum); */
    for (k = 1;k <= rtsNum;k++)
       for (minute = 0;minute <= 2100;minute++)</pre>
           filled[k][minute] = 0;
          /* printf( filled[%d][%d] = %d\n ,k,minute,filled[k][minute]); */
   /* Schedule the overlaps first, before the real supports */
```

```
for (k = 1; k <= rtsWum; k++)
  for (minute = gOverlapBegin[k]; minute <= gOverlapEnd[k]; minute++)</pre>
filled[k][minute] = 1;
  /* Set the Overlaps back to zero; since written to in scheduling 30 Jan */
     for (k = 1; k <= rtsNum; k++)
   gOverlapBegin[k] = 0;
          gOverlapEnd[k] = 0;
    numScheduled = 0;
     /* Schedule each Support in Priority Order */
printf( Spt RTS Beg End Length TAT\n\n );
     for (j = 1;j<= num_supports; j++)</pre>
                                       /* start at 0 for Tsp, 1 for sb */
         spt = order[j];
                        Start SS Order[%d] = %d\n ,j, order[j]); */
            /* printf(
         numScheduled += ScheduleSupport(spt, filled);
          /* end num_supports */
    /* fprintf (stdout, num_supports = %d\n , num_supports); */
   /* The GA wants to MINimize something so do Max poss - actual */
     printf( \n Humber Scheduled = %d\n**********************\n\n\n ,
 numScheduled);
    return(num_supports - numScheduled);
     /* End Eval */
 FUNCTION: read_data
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support data
 * RETURN VALUE: none
 * CALLS: get_coord_array
          read_coords
get_2D_dist_array
          calc_distances
              *****
void
read_data (coord_file, num_supports)
               coord_file[];
char
int
              num_supports;
  int i:
  int spt;
  int rts;
  int bVis, eVis, rLength, tat;
  FILE *fp;
  for (i=1; i<=num_supports; i++)</pre>
gWumWin[i] = 0;
                    /* set number of windows [spt] = 0 */
/* printf( gWumWin[%d]=%d\n ,i,gWumWin[i]); */
  if (fp = fopen(coord_file, r ))
         while (fscanf(fp, %d %d %d %d %d %d , &spt, &rts, &bVis, &eVis, &rLength, &tat ) !=EOF)
         {
              gNumWin[spt] +=1;
              gRTS[spt][gNumWin[spt]] = rts;
              gBeginVis[spt][gNumWin[spt]] = bVis;
```

```
gEndVis[spt][gNumWin[spt]] = eVis;
             gReqLength[spt] = rLength;
             gTurnAroundTime[spt] = tat;
             /* printf ( spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
TAT = %d\n , spt, gRTS[spt][gWumWin[spt]], gBeginVis[spt][gWumWin[spt]],
gEndVis[spt][gHumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
fclose (fp);
/* for (i=1; i<= num_supports;i++)</pre>
printf( gWumWin[%d] = %d\n ,i,gWumWin[i]); */
fatal_error ( Cannot read input city file );
FUNCTION: read_overlap_data
 * DESCRIBE: read in data
 * IMPUT PARAMETERS: filename of list of satellite support overlap data
* RETURN VALUE: none
* CALLS: get_coord_array
         read_coords
         get_2D_dist_array
         calc_distances
********
                          void
read_overlap_data (coord_file)
char
              coord_file[];
 int i;
  int rts;
  int bVis, eVis;
 FILE *fp;
 if (fp = fopen(coord_file, r ))
        while (fscanf(fp, %d %d %d , &rts, &bVis, &eVis ) !=EOF)
             gOverlapBegin[rts] = bVis;
     gOverlapEnd[rts] = eVis;
              /* printf( gOverlapBegin[%d] = %d\t gOverlapEnd[%d] = %d\n ,
rts, gOverlapBegin[rts], rts, gOverlapEnd[rts]); */
             /* printf ( spt = %d, rts= %d, Bvis= %d, Evis = %d, ReqLen = %d,
 TAT = %d\n , spt, gRTS[spt][gHumWin[spt]], gBeginVis[spt][gHumWin[spt]], gEndVis[spt][gHumWin[spt]], gReqLength[spt], gTurnAroundTime[spt] ); */
TAT = %d\n
     /* printf( -----gWumWin[%d] = %d\n , spt,gWumWin[spt]); */
 fclose (fp);
  /* for (i=1; i<= num_supports;i++)</pre>
printf( gWumWin[%d] = %d\n ,i,gWumWin[i]); */
fatal_error ( Cannot read input city file );
                        *****************
FUNCTION: print_overlap_data
 * DESCRIBE: print out data
```

Appendix D. Satellite Range Scheduling Data Processing

The processing of satellite support data begins with raw ASTRO data which is processed by Pascal programs to produce satellite request windows. This data is filtered by a C program to prepare the data for the genetic algorithm.

ASTRO Data

ASTRO (Automated Scheduling Tools for Range Operations) is a computer system and database to aid the range schedulers. The report information for the seven days of test data is in a file named FINLDATA.DFT. This file is the raw input which determines support requests and request visibilities.

Requests and Visibility Windows

Pascal programs written by Gooley and Schalck are used to process the ASTRO data. The programs extract the satellite support requirements and request visibilities. Five programs are used for this algorithm: one to process the low-altitude satellite (low-flyer) support information, and four to process the medium and high-altitude satellite support (high-flyer) information. These programs are run using TurboPascal on a IBM PC compatible computer.

The final data format is:

- 1. Support Number
- 2. RTS the satellite is visible to
- 3. Beginning visibility time of window (in minutes)
- 4. Ending visibility time of window
- 5. Length of support (in minutes)
- 6. TAT (turn around time) required by an RTS
- 7. IRON/Revolution: identifies satellite and pass

LREQ.PAS. This Pascal program reads in the information on low-flyer supports and saves the information to a file called REQLF.DAT containing the low altitude satellite support requests for a day.

HREQ.PAS. This program reads in the information on high-altitude requests for a day and saves the information to two files: 1) REQHF.DAT: the high-altitude requests for a day, and 2) D1V.DAT: the visibilities for the high altitude requests for a day.

TOL.PAS. This program takes tolerance data presented in different formats for each request from HREQ.PAS and standardizes the tolerance window data.

CROSS2.PAS. This program cross checks requests and visibilities. It cross references the visibility file created in HREQ.PAS and the output from TOL.PAS to determine all the RTSs that can satisfy each medium or high altitude support request from TOL.PAS.

RTS.PAS. This Pascal program ensures all RTS sides are included once and only once for each support request-RTS visibility combination.

Prepare Time Window Data for GENITOR

This program reads in the separate data files for the low-flyers and high-flyers and combines them into one file for use in the GENITOR genetic algorithm. It also changes the RTS name to a number for reference in array structures.

```
#include <stdio.h>
int
equal_strings (s1, s2)
char s1[],s2[];
int i = 0, answer;
while ( s1[i] == s2[i] && s1[i] != '\0' && s2[i] != '\0') ++i;
if ( s1[i] == '\0' && s2[i] == '\0')
answer = 1; /* strings equal */
else
answer = 0; /* not equal
return (answer);
int
name_to_Wum (name)
char name[];
int rts;
        (equal_strings (name, "POGO-A") )
if
rts = 1;
else if (equal_strings (name, "POGO-B") )
rts = 2;
else if (equal_strings (name, "POGO-C") )
rts = 3;
else if (equal_strings (name, "HULA-A") )
rts = 4;
else if (equal_strings (name, "HULA-B") )
rts = 5;
else if (equal_strings (name, "COOK-A") )
rts = 6;
else if (equal_strings (name, "COOK-B") )
rts = 7;
else if (equal_strings (name, "IMDI-A") )
rts = 8;
else if (equal_strings (name, "INDI-B") )
rts = 0;
else if (equal_strings (name, "BOSS-A") )
rts = 10;
else if (equal_strings (name, "BOSS-B") )
rts = 11:
else if (equal_strings (name, "LION-A") )
rts = 12:
else if (equal_strings (name,"LIOH-B") )
rts = 13;
else if (equal_strings (name, "GUAM-A") )
rts = 14;
else if (equal_strings (name, "GUAM-B") )
rts = 15:
else if (equal_strings (name, "PIKE-A") )
rts = 16:
else if (equal_strings (name,"PIKE-B") )
rts = 0;
else if (equal_strings (name, "REEF-A") )
rts = 18;
else if (equal_strings (name,"REEF-B") )
rts = 0;
else
printf("UNKNOWN\n");
return(rts);
read_data (out_file, coord_file, start, endPtr)
FILE *out_file;
```

```
char
            coord_file[];
int start;
int *endPtr;
int i; int spt;
   char rtsWame[8];
int bVis, eVis, rLength,tat;
   char iron[16];
   int totWumSpt;
   int rts;
FILE *input_file;
   input_file = fopen(coord_file, "r");
   while (fscanf(input_file, "%d %s %d %d %d %d %s", &spt, &rtsWame, &bVis,
 &eVis, &rLength, &tat, &iron ) !=EOF)
 spt = spt + start:
 totMumSpt = spt;
 rts = name_to_Num(rtsName);
if ( rts != 0 )
            fprintf (out_file,"%d\t%d\t%d\t%d\t%d\t%d\t%d\n", spt, rts, bVis, eVis,
 rLength, tat);
   fclose (input_file);
   *endPtr = totWumSpt;
/****
           MAIN PROGRAM ****/
main()
{ /* start MAIN */
/**** Initialize and declare Variables *****/
char lfdata[12], hfdata[12];
int WumLFSpt, WumHFSpt;
FILE *output_file;
output_file = fopen("datafile", "w");
/**** Input LF Filename *****/
printf("What is the LF filename?");
scanf ("%s", lfdata);
/**** Read in LF Data and print to Dayout.dat *****/
read_data (output_file, lfdata, 0, &WumLFSpt);
/**** Input HF filename *****/
printf("What is the HF filename?");
scanf ("%s", hfdata);
/**** Read in HF Data and print to Dayout.dat *****/
read_data (output_file, hfdata, MumLFSpt, & MumHFSpt);
fclose (output_file);
printf("MumLFSpt = %d, EndHFSpt = %d\n", MumLFSpt, MumHFSpt);
} /* end Main */
```

Appendix E. Schedules for Week One Data

Schedules produced for each of seven days of data follow. The scheduled supports are sorted by RTS, and by time. The columns are: Support number, RTS, Beginning Time, Ending Time, Support Length, and Support turn around time.

255 POGO-A 890 890 21 1 20 277 HULA-A 996 1 1 20 1 15 276 POGO-A 971 987 16 20 277 HULA-A 996 10 16 20 1 15 276 POGO-A 1020 1035 15 15 15 112 HULA-A 1079 1095 16 20 1 280 POGO-A 1089 1099 10 15 296 HULA-A 1110 1120 10 15 15 16 POGO-A 1089 1099 10 15 296 HULA-A 1150 1120 10 15 11 12 POGO-A 1171 1188 17 20 296 HULA-A 1120 1285 25 15 127 POGO-A 1221 1236 12 20 21 15 15 124 POGO-A 1221 1236 12 20 21 15 15 124 POGO-A 1260 1276 14 20 21 15 HULA-A 1325 1335 10 15 127 POGO-A 1260 1276 14 20 21 15 HULA-A 1325 1335 10 15 134 POGO-A 1260 1276 14 14 20 21 15 15 15 15 15 15 15 15 15 15 15 15 15	Day 1	142 POGO-C 1335 1351 16 20 154 HULA-A 0 35 35 15
30 P0G0-C 326 342 16 20 294 C00K-A 1145 1160 15 15 37 P0G0-C 364 376 12 20 126 C00K-A 1192 1206 14 20 45 P0G0-C 449 466 17 20 138 C00K-A 1293 1308 15 20 11 P0G0-C 505 515 10 15 316 C00K-A 1365 1370 5_15	245 POGO-A 870 890 20 15 94 POGO-A 939 950 11 20 99 POGO-A 971 987 16 20 276 POGO-A 1020 1035 15 15 109 POGO-A 1062 1074 12 20 280 POGO-A 1089 1099 10 15 116 POGO-A 1171 1188 17 20 124 POGO-A 1171 1188 17 20 124 POGO-A 1221 1235 14 20 134 POGO-A 1221 1235 14 20 139 POGO-A 1318 1333 15 20 145 POGO-A 1318 1333 15 20 145 POGO-A 1415 1450 35 15 158 POGO-B 1415 1450 35 15 158 POGO-B 51 5 10 15 4 POGO-B 51 64 13 20 172 POGO-B 105 10 15 15 15 POGO-B 169 179 10 20 21 POGO-B 169 179 10 20 21 POGO-B 375 390 15 15 42 POGO-B 426 443 17 20 200 POGO-B 376 390 15 15 42 POGO-B 4643 17 20 203 POGO-B 475 520 45 15 56 POGO-B 656 639 14 20 67 POGO-B 656 639 14 20 67 POGO-B 666 67 14 20 236 POGO-B 834 848 14 20 252 POGO-B 834 848 14 20 252 POGO-B 1047 1061 15 267 POGO-B 1047 1061 14 20 284 POGO-B 1047 1061 14 20 284 POGO-B 1047 1061 14 20 284 POGO-B 1047 1061 15 15 277 POGO-B 1047 1061 15 15 277 POGO-B 1047 1061 15 15 279 POGO-B 1145 1188 13 20 296 POGO-B 1147 1262 15 20 296 POGO-B 1147 1262 15 20 296 POGO-B 1148 1128 1208 10 15 131 POGO-B 1247 1262 15 20 200 POGO-B 1247 1262 15 20	264 HULA-A 932 947 15 15 98 HULA-A 970 981 11 20 15 17 HULA-A 1035 1040 5 16 11 12 HULA-A 1035 1040 5 16 11 12 HULA-A 1079 1095 16 20 15 11 12 HULA-A 1170 1120 10 15 15 15 15 15 15 15 15 15 15 15 15 15
\text{\text{P} \text{\text{P} \text{\text{P} \text{\text{\text{P} \text{\tint{\text{\text{\text{\text{\text{\text{\text{\text{\text{\tin\text{\texi}\text{\text{\texi}\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\text{\t	119 POGO-B 1145 1158 13 20 296 POGO-B 1173 1183 10 15 293 POGO-B 1198 1208 10 15 131 POGO-B 1247 1262 15 20 308 POGO-B 1290 1330 40 15 146 POGO-B 1362 1375 13 20 168 POGO-C 70 105 35 15	44 COOK-A 441 452 11 20 216 COOK-A 525 535 10 15 219 COOK-A 550 575 25 15 231 COOK-A 660 705 45 15 242 COOK-A 735 765 30 15 242 COOK-A 825 840 15 15 88 COOK-A 874 890 16 20 261 COOK-A 905 920 15 15 100 COOK-A 975 989 14 20 282 COOK-A 1080 1095 15 15 283 COOK-A 1110 1130 20 15

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213 LION-A 515 535 20 15
68 LION-A 665 678 13 20
233 LION-A 693 713 20 15
84 LION-A 846 857 12 20
256 LION-A 940 955 15 20
269 LION-A 990 1010 20 15
106 LION-A 1039 1052 13 20
117 LION-A 1252 1267 15 20
314 LION-B 225 239 14 20
28 LION-B 225 239 10 35 15
20 LION-B 320 330 10 15
210 LION-B 320 330 10 15
210 LION-B 860 670 10 15
87 LION-B 860 670 10 15
87 LION-B 933 947 14 20
91 LION-B 933 947 14 20
101 LION-B 1052 1063 11 20
1120 LION-B 1052 1063 11 20
1120 LION-B 1245 1270 25 15
8 GUAM-A 108 123 15 20
118 GUAM-A 108 123 15 20
120 LION-B 1245 1270 25 15
8 GUAM-A 265 285 20 15
189 GUAM-A 266 651 15 20
207 GUAM-A 355 369 14 20
207 GUAM-A 355 550 15 20
64 GUAM-A 366 651 15 20
93 GUAM-A 366 651 15 20
105 GUAM-A 108 1112 14 20
105 GUAM-A 108 1131 13 20
114 GUAM-A 108 1112 14 20
105 GUAM-A 1283 1299 16 20
313 GUAM-A 1350 1370 20 15
16 GUAM-B 173 187 14 20
107 GUAM-B 1360 1150 10 15
147 GUAM-B 1360 1370 20 15
149 GUAM-A 108 112 14 20
105 GUAM-A 108 112 14 20
112 FIKE-A 138 164 16 20
23 PIKE-A 440 454 14 20
15 PIKE-A 430 454 14 20
15 PIKE-A 400 454 14 20
15 PIKE-A 358 415 57 15
43 PIKE-A 438 154 16 20
20 1 PIKE-A 614 629 15 20
18 REEF-A 193 12 10
18 REEF-A 193 12 10
18 REEF-A 193 12 10
18 REEF-A 193 20 7 15
18 REEF-A 193 12 20
18 REEF-A 480 496 16 20
59 REEF-A 632 644 12 20
72 REEF-A 893 946 13 20
129 REEF-A 893 946 13 20
129 REEF-A 893 946 13 20
129 REEF-A 1590 1445 55 15
Day 2
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           145 POGO-A 45 90 45 15 15 152 POGO-A 105 120 15 15 15 157 POGO-A 150 170 20 15 162 POGO-A 200 215 15 15 166 POGO-A 295 305 10 15 173 POGO-A 295 305 10 15 180 POGO-A 419 435 16 20 36 POGO-A 460 475 15 20 41 POGO-A 512 529 17 20 45 POGO-A 554 570 16 20 52 POGO-A 655 660 15 15 61 POGO-A 656 660 15 15 61 POGO-A 690 705 15 20 222 POGO-A 750 780 30 15 73 POGO-A 808 823 15 20 78 POGO-A 858 868 10 20 85 POGO-A 960 975 15 15
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94 POGO-A 1015 1027 12 20 262 POGO-A 1085 1075 10 15 265 POGO-A 10805 1075 10 15 103 POGO-A 11302 1147 15 20 212 POGO-A 11302 1147 15 20 220 POGO-A 11302 1147 15 15 113 POGO-A 11302 1246 14 20 2130 POGO-A 1232 1246 15 15 1137 POGO-A 1231 1346 15 20 2124 POGO-A 1331 1346 15 20 2128 POGO-A 1231 1447 16 20 138 POGO-A 1431 1447 16 20 138 POGO-B 0 10 10 15 15 149 POGO-B 65 100 35 15 15 149 POGO-B 135 150 15 15 156 POGO-B 135 150 15 15 156 POGO-B 250 255 5 15 15 156 POGO-B 270 255 5 15 15 1570 POGO-B 270 255 15 15 1570 POGO-B 360 375 15 15 181 POGO-B 360 375 15 15 181 POGO-B 360 375 15 15 181 POGO-B 360 375 15 15 186 POGO-B 360 375 15 15 187 POGO-B 482 497 15 15 187 POGO-B 1304 1040 16 197 POGO-B 1304 1040 10 157 POGO-B 1304 1040 10 120 POGO-B 1304 1040 10 120 POGO-B 1304 1040 10 121 POGO-B 1203 104 10 122 POGO-B 1304 1040 10 123 POGO-B 1304 1040 10 129 POGO-B 1304 1040 10 139 POGO-B 1304 1040 10 139 POGO-B 1304 1040 10 139 POGO-C 210 240 30 15 184 POGO-B 1304 1040 10 184 POGO-C 230 300 10 15 185 POGO-C 2470 480 10 15 189 POGO-C 2470 480 10 15 189 POGO-C 256 672 16 20 175 POGO-C 656 672 16 20 175 POGO-C 656 672 16 20 175 POGO-C 657 672 16 20 175 POGO-C 1058 1076 17 20 188 POGO-C 1058 1076 17 20 188 POGO-C 1058 1076 17 20 189 POGO-C 1058 1076 10 15 189 POGO-C 1131 1143 1123 112 20 POGO-C 1131 1143 112 20 21 POGO-C 1206 1226 16 20 21 293 POGO-C 1351 1351 1366 15 10 POGO-C 1365 1076 10 10 15 10 POGO-C 136	248 HULA—A 945 950 5 15 249 HULA—A 965 975 10 15 254 HULA—A 960 1010 20 15 97 HULA—A 1050 1067 17 20 100 HULA—A 1101 1116 18 20 269 HULA—A 1200 1215 30 15 278 HULA—A 1200 1215 35 15 278 HULA—A 1305 1315 10 15 278 HULA—A 1305 1315 10 15 290 HULA—A 1305 1315 10 15 290 HULA—A 1300 1365 15 15 291 HULA—B 160 165 5 15 291 HULA—B 160 165 5 15 291 HULA—B 160 165 5 15 291 HULA—B 482 497 15 20 292 HULA—B 482 497 15 20 209 HULA—B 482 497 15 20 217 HULA—B 935 950 15 15 21 HULA—B 935 950 15 15 21 HULA—B 1305 105 102 17 207 HULA—B 1305 1363 13 246 HULA—B 1180 1285 1296 10 15 273 HULA—B 1350 1363 13 15 140 COOK—A 282 297 15 20 24 COOK—A 282 297 15 20 24 COOK—A 282 297 15 20 24 COOK—A 470 495 25 15 228 COOK—A 282 297 120 120 190 15 228 COOK—A 795 830 35 15 238 COOK—A 1624 1375 115 228 COOK—A 1624 1375 15 15 228 COOK—A 1244 1260 16 20 248 COOK—A 1244 1260 16 20 258 COOK—A 1244 1270 120 190 15 226 COOK—B 120 120 190 15 226 COOK—B 120 120 190 15 226 COOK—B 120 120 190 15 226 COOK—B 1306 115 15 227 COOK—B 120 120 190 15 228 COOK—B 1364 1375 15 229 COOK—B 1360 115 15 229 COOK—B 1360 115 15 229 COOK—B 1360 110 15 229 COOK—B 1360 110 15 229 COOK—B 1360 110 15 229 COOK—B 1360 10 15 220 COOK—B 1360 10 15 221 110 1—A 450 10 10 10 15 221 110 1—A 450 10 10 10 15 223 110 1—A 450 10 10 10 15 234 110 1—A 340 345 355 10 15 239 110 1—A 360 60 60 5 5 15 240 110 1—A 360 60 60 5 5 15 241 110 1—A 360 60 60 5 5 15 242 110 1—A 360 60 60 5 5 15 243 110 1—A 360 60 60 5 5 15 244 110 1—A 360 60 60 5 5 15 245 110 1—A 360 60 60 5 5 15 246 110 1—A 360 60 60 5 5 15 257 110 1—A 360 60 60 5 5 15 258 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A 360 60 60 5 5 15 259 110 1—A
151 HULA-A 150 160 10 15 161 HULA-A 180 190 10 15 15 HULA-A 225 241 16 20 167 HULA-A 325 276 20 15 23 HULA-A 331 347 16 20 27 HULA-A 331 347 16 20 183 HULA-A 383 398 15 20 183 HULA-A 413 418 5 15 191 HULA-A 435 450 15 15 195 HULA-A 480 490 10 15 195 HULA-A 510 530 20 15 202 HULA-A 545 550 5 15 202 HULA-A 546 550 5 15 202 HULA-A 660 670 10 15 64 HULA-A 706 722 16 20 224 HULA-A 706 722 16 20 224 HULA-A 775 780 5 15 226 HULA-A 795 800 5 15 232 HULA-A 815 820 5 15 238 HULA-A 865 885 20 15 84 HULA-A 907 920 13 20	234 INDI-A 835 840 5 15 239 INDI-A 870 890 20 15 241 INDI-A 905 915 10 15 247 INDI-A 935 955 20 15 93 INDI-A 1007 1019 12 20

165 BOSS-A 225 230 5 15 174 BOSS-A 270 290 20 15 22 BOSS-A 310 324 14 20 30 BOSS-A 452 466 14 20 198 BOSS-A 452 466 14 20 198 BOSS-A 560 570 10 15 205 BOSS-A 664 621 17 20 55 BOSS-A 664 621 17 20 55 BOSS-A 664 621 17 20 206 BOSS-A 671 691 20 15 214 BOSS-A 706 716 10 15 68 BOSS-A 706 716 10 15 68 BOSS-A 773 793 20 15 214 BOSS-A 790 910 10 15 87 BOSS-A 940 951 11 20 250 BOSS-A 940 951 15 15 96 BOSS-A 1040 1056 16 20 107 BOSS-A 1040 1056 16 20 107 BOSS-A 1195 1200 5 15 282 BOSS-A 1195 1200 15 15 282 BOSS-A 1195 1200 15 15 297 BOSS-A 1195 1200 15 15 297 BOSS-A 1380 1400 20 15 135 BOSS-A 1380 1400 20 15 135 BOSS-B 60 360 300 15	50 PIKE-A 605 617 12 20 56 PIKE-A 648 660 14 20 63 PIKE-A 701 716 15 20 77 PIKE-A 844 860 16 20 105 PIKE-A 1225 1240 15 20 112 PIKE-A 1266 1276 10 20 118 PIKE-A 1360 1375 15 20 128 PIKE-A 1360 1375 15 20 296 PIKE-A 1390 1410 20 15 10 REEF-A 151 167 16 20 33 REEF-A 451 465 14 20 200 REEF-A 515 535 20 15 47 REEF-A 591 602 11 20 58 REEF-A 661 676 15 20 72 REEF-A 803 818 15 20 79 REEF-A 859 872 13 20 109 REEF-A 1301 1217 16 20 122 REEF-A 1322 1334 12 20
48 BOSS-B 602 616 14 20 208 BOSS-B 631 636 5 15 216 BOSS-B 675 680 5 15	Day 3
165 BOSS-A 226 230 5 15 174 BOSS-A 270 290 20 15 20 BOSS-A 310 324 14 20 34 BOSS-A 410 426 16 20 34 BOSS-A 452 466 14 20 198 BOSS-A 560 570 10 15 205 BOSS-A 560 570 10 15 205 BOSS-A 664 621 17 20 55 BOSS-A 664 621 17 20 55 BOSS-A 664 621 17 20 55 BOSS-A 664 621 17 20 56 BOSS-A 664 621 17 20 57 BOSS-A 664 621 17 20 58 BOSS-A 664 626 15 20 206 BOSS-A 664 626 15 20 214 BOSS-A 706 716 10 15 68 BOSS-A 706 716 10 15 68 BOSS-A 406 981 11 20 225 BOSS-A 960 981 15 20 242 BOSS-A 960 981 15 20 242 BOSS-A 960 981 15 15 87 BOSS-A 1040 1056 16 20 107 BOSS-A 1146 1177 12 20 250 BOSS-A 1040 1056 16 20 107 BOSS-A 1146 1177 12 20 287 BOSS-A 122 1438 16 20 287 BOSS-A 126 16 16 20 274 BOSS-A 126 16 16 20 274 BOSS-A 126 16 16 20 274 BOSS-A 126 16 16 20 287 BOSS-B 602 616 14 20 280 BOSS-B 602 616 14 20 280 BOSS-B 602 616 14 20 208 BOSS-B 602 616 14 20 208 BOSS-B 602 616 14 20 208 BOSS-B 1040 1050 15 15 260 BOSS-B 1040 1050 15 15 260 BOSS-B 1040 1050 10 15 281 BOSS-B 1270 1285 215 281 BOSS-B 1320 1285 165 153 LION-A 110 120 10 15 133 LION-A 209 222 13 20 188 LION-A 336 351 16 20 288 LION-A 437 453 16 20 288 LION-A 1209 1219 10 15 287 LION-A 1306 1219 10 15 287 LION-A 1306 1219 10 15 288 LION-A 1306 122 13 20 288 LION-A 1307 985 14 20 290 LION-A 1308 1219 10 15 201 LION-A 1308 1219 10 15 201 LION-A 1308 1219 10 15 201 LION-B 526 550 250 201 LION-B 656 670 15 16 201 LION-B 656 670 15 16 201 LION-B 656 670 15 16 201 LION-B 735 751 16 20 201 LION-B 1063 1075 12 20 201 LION-B 868 880 12 20 202 110 B 1304 1305 10 20 203 LION-B 1063 1075 12 20 204 LION-B 1063 1075 12 10 205 LION-B 1063 1075 12 10 206 LION-B 1063 1075 12 10 207 GUAM-A 1307 129 10 10 15 208 LION-B 1063 1075 12 10 209 GUAM-A 1308 1306 15 20 209 GUAM-A 1308 1306 15 20 201 LION-B 1063 1075 12 10 201 LION-B 1063 1075 12 10 201 LION-B 1066 60 60 203 GUAM-A 1000 1066 60 204 PIKE-A 509 525 16 20 205 GUAM-A 1000 1660 20 206 GUAM-A 1000	166 POGO-A 1 1 0 0 152 POGO-A 30 60 30 15 8 POGO-A 92 108 16 20 164 POGO-A 135 150 15 15 180 POGO-A 240 255 15 15 125 POGO-A 288 304 16 20 130 POGO-A 380 375 15 15 137 POGO-A 380 375 15 15 137 POGO-A 390 342 12 20 135 POGO-A 390 342 12 20 135 POGO-A 490 506 16 20 215 POGO-A 490 506 16 20 215 POGO-A 490 506 16 20 215 POGO-A 570 586 11 20 42 POGO-A 676 689 13 20 72 POGO-A 570 770 20 15 79 POGO-A 804 818 14 20 234 POGO-A 750 770 20 15 79 POGO-A 804 818 14 20 238 POGO-A 833 838 5 15 244 POGO-A 899 915 16 20 238 POGO-A 899 915 16 20 238 POGO-A 930 945 15 15 257 POGO-A 116 116 1128 12 20 251 POGO-A 900 945 15 15 257 POGO-A 116 1128 12 20 251 POGO-A 116 116 1128 12 20 277 POGO-A 1174 1190 16 20 283 POGO-A 1205 1215 10 15 289 POGO-A 1205 1215 10 15 289 POGO-A 1273 1289 16 20 171 POGO-A 1174 1190 16 20 283 POGO-A 1273 1289 16 20 172 POGO-B 1174 120 15 15 127 POGO-B 1174 120 15 15 127 POGO-B 1273 1289 16 20 133 POGO-B 1320 1335 15 20 170 POGO-B 1405 140 35 15 15 POGO-B 480 16 20 159 POGO-B 851 10 15 15 POGO-B 15 15 15 12 POGO-B 165 180 15 20 170 POGO-B 428 441 13 20 14 POGO-B 448 456 461 10 15 15 POGO-B 448 456 461 10 15 16 POGO-B 498 514 16 20 170 POGO-B 456 461 10 15 16 POGO-B 498 514 16 20 170 POGO-B 498 514 16 20 170 POGO-B 498 514 16 20 170 POGO-B 498 514 16 20 180 180 180 180 180 100 100 100 100 100

113 121 290 128 139 151 7 P(163 17 I 186 196	P0G0-B P0G0-B P0G0-B P0G0-B P0G0-C P0G0-C P0G0-C P0G0-C	1146 1218 1247 1289 1350 30 103 120 189 20 285 360	1162 1232 1237 1357 1366 135 135 14 135 14 135 14 135 14 15 15 15 15 15 15 16 16 17 16 17 16 17 16 17 16 17 16 17 16 17 17 17 17 17 17 17 17 17 17 17 17 17	16 20 14 20 10 120 14 20 16 20 5 15 20 15 15
205 I 47 I 216 I 57 I 248 I 2269 1279 2285 131	BBBBC9C CCCC C CCCCCCCCCCCCCCCCCCCCCCCC	420 4 501 5 599 6 710 7 797 8 900 9 1060 1155 1201	135 15 15 16 15 16 13 16 13 16 145 1070 11117 11270 11314	15 20 15 20 15 10 15 10 10 10 10 10 10 10 10 10 10 10 10 10
138 148 158 108 1181 1185 1194 1383	POGO-C HULA-A HULA-A HULA-A HULA-A HULA-A HULA-A HULA-A HULA-A HULA-A	1347 30 91 122 5 212 6 212 2 280 315 340 398	1360 15 15 15 16 16 17 18 17 18 19 10 10 10 10 10 10 10 10 10 10	13 20 55 20 15 20 15 15 15 15 15 20
2107 2217 6231 2237 2249 256 271 271 2112	HULA-A	495 676 69 710 780 869 8 915 935 1022 1071 1124	570 55 590 15 720 16 720 5 84 15 940 5 940 5 940 38 1086 1137	15 15 15 15 15 15 15 15 15 15 16 16 13 15 13 13 13 13 13 13 13 13 13 13 13 13 13
287 288 298 203 147 155 179 188	HULA-A HULA-A HULA-A HULA-A HULA-A HULA-B HULA-B HULA-B HULA-B HULA-B HULA-B HULA-B HULA-B	1170 1230 1260 1295 1345 0 10 145 195 240 379	1200 1240 1275 1295 1330 10 15 10 165 20 165 2155 16	30 15 15 15 15 15 15 15 15 15 15 15 15 15 15 15 1
2230 76 96 102 275 281 124 298	HULA-B HULA-B HULA-B HULA-B HULA-E HULA-E HULA-E	565 705 778 7 991 1 1034 1105 1175 1245 0 374 0 375 315 3	805 40 750 48 92 14 007 16 1048 1120 11258 1258 14 20 220 45	15 20 14 20 14 20 15 15 15 13 20 40 15
197 204 2011 2014 2014 2229 239 85	COOK - A COOK - A	175 315 375 420 495 520 580 690 790 858 917	30 15 390 15 435 15 505 10 525 5 565 25 590 10 705 15 874 16 33 16	20 15 15 15 15 15 15 15 15 20 20

141 BOSS-B 1369 1379 10 20 308 BOSS-B 1394 1414 20 15 161 LIOW-A 85 105 20 15 18 LIOW-A 239 253 14 20 27 LIOW-A 338 32253 14 20 27 LIOW-A 337 347 10 15 198 LIOW-A 375 380 5 15 203 LIOW-A 420 480 60 15 428 LIOW-A 690 710 20 15 80 LIOW-A 690 710 20 15 80 LIOW-A 690 710 20 15 80 LIOW-A 690 710 11 14 20 918 LIOW-A 997 1011 14 20 918 LIOW-A 997 1011 14 20 118 LIOW-A 1094 1108 14 20 118 LIOW-A 1094 1108 14 20 118 LIOW-A 1393 1402 9 20 118 LIOW-A 1393 1402 9 20 118 LIOW-B 60 180 120 15 157 LIOW-B 60 180 120 15 157 LIOW-B 800 845 35 15 90 LIOW-B 336 350 14 20 200 LIOW-B 390 420 30 15 241 LIOW-B 810 845 35 15 90 LIOW-B 300 913 13 20 99 LIOW-B 1080 1100 20 15 310 LIOW-B 1080 1100 20 15 310 LIOW-B 1080 1100 20 15 310 LIOW-B 300 913 13 20 99 LIOW-B 300 913 13 20 99 LIOW-B 1080 1100 20 15 305 LIOW-B 300 1310 10 15 15 GUAM-A 171 183 12 20 272 LIOW-B 1080 1100 20 15 305 LIOW-B 1080 1100 20 15 305 LIOW-B 1080 1100 20 15 305 LIOW-B 1080 1100 20 15 306 LIOW-B 1080 1100 20 15 307 GUAM-A 797 310 13 20 94 GUAM-A 212 229 17 20 28 GUAM-A 203 217 120 39 GUAM-A 1064 1078 14 20 274 GUAM-A 797 310 13 20 94 GUAM-A 1064 1078 14 20 275 LIOW-B 1080 1108 120 95 155 GUAM-A 1295 1305 10 15 135 GUAM-A 1295 1305 10 15 135 GUAM-A 1296 1306 10 15 136 GUAM-A 1296 1306 10 15 137 LIOW-B 1300 1310 10 15 137 LIOW-B 1300 1310 10 15 15 GUAM-A 1064 1078 14 20 107 PIKE-A 180 360 180 10 15 137 LIOW-B 1300 130 10 15 138 GUAM-B 1036 1046 10 20 172 PIKE-A 180 360 180 15 154 PIKE-A 635 648 13 20 172 PIKE-A 180 360 180 15 154 PIKE-A 635 648 13 20 172 PIKE-A 180 360 180 15 110 PIKE-A 635 648 13 20 172 PIKE-A 180 360 180 15 110 PIKE-A 180 180 180 180 180 180 180 180 180 180	
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Day 4

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Vita

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